## Homework #1

## Math 231b

## "Due": February 15th, 2017

Guidelines:

- Type up your solution to the assignment in LATEX. (You might want to avail yourself of the excellent diagrams package tikz-cd.)
- Submit the PDF via Canvas, in the Assignments section.

Failure to meet these guidelines may result in loss of points.

**Problem 1.** Consider the unreduced mapping cylinder  $M_f$  of a map  $f: X \to Y$ , together with its inclusion  $i: Y \to M_f$ . Show that the usual retraction  $r: M_f \to Y$  of i (which projects down from the interval coordinate) induces an isomorphism  $r_*: \pi_*M_f \to \pi_*Y$ . (We used this fact when constructing the homotopy long exact sequence of an arbitrary pair.)

**Problem 2.** Show that if  $f: X \to Y$  is a homotopy equivalence, then  $f_*: \pi_*(X, x_0) \to \pi_*(Y, f(x_0))$  is an isomorphism for all choices of  $x_0 \in X$  and  $* \ge 0$ .

**Problem 3.** In class, I used the mysterious phrase "limit condition" twice. Given a functor  $F: J \to C$ , thought of as a J-shaped diagram in C, we define a *cone* of F to be a constant functor  $x: J \to C$  together with a natural transformation  $x \to F$ . A *limit* of F is a terminal object in the category of cones.

1. Expand the definition of natural transformation and constant functor to reveal that a cone is equivalent to the data of an object  $x \in C$  together with maps  $f_j: x \to F(j)$  for each object  $j \in J$  such that for any map  $g: j \to j'$  in the diagram there is a commuting triangle



- 2. Now expand the definition of limit to see that a limit, expressed as an object  $\ell$  together with maps  $h_j$ , has the property that any other cone point x and its maps  $f_j$  factor uniquely through a map  $x \to \ell$ .
- 3. Show that the product  $X \times Y$  is the limit of the diagram with objects X and Y and no non-identity arrows.
- 4. Show that the equalizer E of a pair of functions  $X \rightrightarrows Y$  of sets is indeed the limit of that diagram.

**Problem 4.** A functor  $G: \mathbb{C}^{\mathrm{op}} \to \mathsf{Sets}$  is called *representable* when there exists an object Y and a natural isomorphism

$$G \xrightarrow{\simeq} \mathsf{Sets}(-, Y).$$

1. From a morphism  $t: Y \to Y'$  of representing objects, construct a natural transformation  $t_*: G \to G'$  of the functors they represent.

- 2. From a natural transformation  $G \to G'$  of represented functors, construct a morphism  $Y \to Y'$  of the representing objects.
- 3. Show also that your assignments respect composition of natural transformations and of morphisms.
- 4. Show that your assignments are mutual inverses, i.e., a natural transformation of representable functors is exactly the same information as a morphism of representing objects.

Congratulations! You have proved the Yoneda lemma: the functor

$$C \rightarrow Categories(C^{op}, Sets)$$

describes a fully faithful embedding.

**Problem 5.** Explain convincingly why the usual recipe for forming a group structure on  $\pi_1(X, x_0)$  does not apply to the set of relative homotopy classes  $\pi_1(I, \partial I)$ .

**Problem 6.** Let  $p: E \to B$  be a map and consider

$$Z = \{(e, \gamma) \in E \times B^I \colon p(e) = \gamma(0)\} \subseteq E \times B^I.$$

A path lifting function for p is a map  $\lambda: Z \to E^I$  with  $\lambda(e, \gamma)(0) = e$  and  $p \circ \lambda(e, \gamma) = \gamma$ .

- 1. Show that p is a fibration if and only if there is a path lifting function  $\lambda$  for p.
- 2. Let  $p: E \to B$  be a fibration with fiber F, and let  $P_p$  be the pathspace construction p described in class. Given a path lifting function  $\lambda: Z \to E^I$  for p, define maps

$$g: F \to P_p, \qquad f \mapsto (f, \omega_0),$$
  

$$h: P_p \to F, \qquad (e, \gamma) \mapsto [\lambda(e, \gamma^{-1})](1),$$

where " $\gamma^{-1}$ " denotes the path  $\gamma$  run backwards. Show that g and h present the two halves of a homotopy equivalence.