

Introduction/Summary

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Office hours: 321h SC, T1-2pm, W11-12pm.

CAs: Thayer Anderson, Davis Lazarecki, Haubong Park, Rohil Prasad.

Grades: Homework due Wednesday morning, before class begins, separated by CA. (25%) ~~8~~ **LaTeX**

• A midterm: 10/26, in class (25%).

• Final exam: 12/10, 9am (50%).

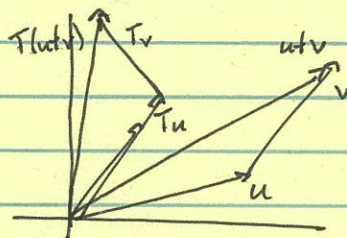
CAs have office hours. Weekly problem night: 11:30pm, Leverett.

A function is "linear" if $T(c \cdot x) = c \cdot T(x)$ and $T(x+y) = T(x) + T(y)$.

Ex: The only such $f^{\text{us}} \mathbb{R} \rightarrow \mathbb{R}$ are $T(x) = k \cdot x$ for some $k \in \mathbb{R}$.

↳ But there are more with other domains + codomains.

Exs: A rotation of \mathbb{R}^2 :



Evaluation of polynomials:
 $T(f) = f(1)$.

Derivatives: $\frac{d}{dx}(f(x) + g(x)) = \frac{df}{dx} + \frac{dg}{dx}$, and $\frac{d}{dx}(c \cdot f(x)) = c \cdot \frac{df}{dx}$.

Linear algebra is about studying these T 's + the equations

they appear in: $T(x) = y$ ← how can we solve this for fixed y ?

$T(x) = x$ ← how about this, with x on both sides?

$T(u) \in U$

Gaussian elimination

① + matrix representations,
basic structure of vector spaces

eigenvectors and eigenvalues, ②
the Spectral Theorem + SVD,

generalized eigenvalues, quadratic f^{us} ,
Jordan normal form. ③

Main goals for this class:

Linear algebra in its own right: tangible, successful mathematics.

Linear algebra for math. apps.: calculus next semester.

Proof-writing. Manipulating def^{us}. Math as simulation + substrate.

Proof techniques I

Major goal of this class: learning to write proofs. Proofs divide into two main camps: algebraic + analytic. There are a/b. To start, we will consider some basic, universal techniques.

Mathematics is about designing models + then arguing about their behaviors. It is important to become a good + flexible debater, and to be eager to consider all points of view — mathematics is very rigid, but mathematicians are highly fallible.

Ex: A ^{whole #} ~~number~~ ^{integer} is divisible by 9 exactly if the sum of its decimal digits is.

First, examples: $81 = 9 \cdot 9$, and $8+1=9=9 \cdot 1$.

$693 = 9 \cdot 77$, and $6+9+3=18=9 \cdot 2$.

Meanwhile, $500 = 4 \cdot 5^3$, and $5+0+0=5$.

convert
to
symbols.

not allowed to pick an example

Pf: Suppose that n is ^{a pos.} integer. Its decimal expansion is

$$n = 10^k a_k + 10^{k-1} a_{k-1} + \dots + 10 a_1 + a_0 \text{ for } 0 \leq a_i \leq 10.$$

The digital sum is then

$$s(n) = a_k + a_{k-1} + \dots + a_1 + a_0.$$

There are very similar expressions, so — on a lark — we subtract them:

$$\Delta := n - s(n) = (10^k - 1)a_k + (10^{k-1} - 1)a_{k-1} + \dots + (10 - 1)a_1 + (1 - 1)a_0.$$

This is divisible by 9, since each summand has a factor like $99 \dots 99$, ^{no matter what n was}

We then have $n = \Delta + s(n)$ and $s(n) = n - \Delta$. So if

$s(n)$ is div. by 9 as

n is div. by 9

then

n is div. by 9

$s(n)$ is div. by 9

because the sum/difference of two things divisible by 9 is again so. \square

Thought experiment (Wason selection task):

You are told: every card has a number on one side and a color on the other. If the number is even, then the color must be red.



x ✓ ✓ x

How many cards do you need to check to verify the claim (underlined)?

~~Drunk driving is illegal.~~ Underage drinking is illegal. If you're drunk, you must be ≥ 21 .

28

14

Drunk

Sobor

This is meant to illustrate the contrapositive:

logically equivalent statements $\left\{ \begin{array}{l} \text{If you are drunk, then you must be } \geq 21. \\ \text{If you are not } \geq 21, \text{ then you must not be drunk.} \end{array} \right.$

However, sometimes one is easier than the other.

~~Ex: Show that there exist two irrational #'s with a rational.~~

Ex: Show that if xy and $x+y$ are even, then x and y are even.

HERE IS THE CONTRAPOSITIVE

Pf: We instead show that if x and y are not both even, then xy and $x+y$ will not both be even.

Case 1: x even, y odd means $x+y$ is odd.

Case 2: x odd, y even means $x+y$ is odd.

Formally identical, since $xy = yx$. This is

sometimes phrased as "Assume one of the two is odd and the other even. Without loss, we may take x odd & y even."

Case 3: x and y both odd means $x \cdot y$ is odd.

Interstitial examples:

Also do a direct version of this.
+ \Rightarrow same parity

x	y	$x+y$	xy
2	1	3	2
3	4	7	12
5	7	12	35
8	6	14	48

odd!
even!

Another powerful proving tool for statements indexed by \mathbb{N} is induction.

Ex: $1 + 4 + 9 + \dots + n^2 = \frac{n(n+1)(2n+1)}{6}$ for all n .

Pf: For $n=1$, $1 = 1 \cdot (1+1) \cdot (2 \cdot 1 + 1) / 6$. \checkmark

Assume $1 + 4 + 9 + \dots + j^2 = \frac{j(j+1)(2j+1)}{6}$ for some j .

Then $1 + 4 + 9 + \dots + j^2 + (j+1)^2 = \frac{j(j+1)(2j+1)}{6} + (j+1)^2$

$= \frac{(j+1)}{6} (j(2j+1) + 6(j+1)) = \frac{(j+1)}{6} (2j^2 + 2j + 6j + 6) = \frac{(j+1)(j+2)(2j+3)}{6} \checkmark$

Yihang: Fri 11am-12pm

Proof techniques II

Today we talk about two more complicated aspect of proofwriting: quantification and contradiction.

We saw quantification yesterday: when we showed that all integers n were divisible by 9 exactly when their digital sum are, the "all" is a quantifier. There is a second kind of quantifier, also of interest:

There are claims about general behavior.

"There is a solution x to the equation $x^2 + x = 0$."

called an existential quantifier. There are claims about examples, (Pf: Pick $x = 0$ or $x = -1$.) and they are often short.

These are interrelated:

If not all x satisfy P , then there must exist an x not satisfying P .

If there does not exist an x satisfying P , then all x must not satisfy P .

Moving "not" past the quantifier changes it! This is the font of proof by counterexample: if you want to show that not all x have property P , then you need exhibit only one such x .

Ex: Falsify the statement that for any $y \in \mathbb{R}$ there is an $x \in \mathbb{R}$ with $y = x^2$.

Pf: We need to show that there is a y which for any x , $y \neq x^2$.

If we select $y = -1$, then any x has x^2 nonnegative, hence $y \neq x^2$. \square

It is also possible to prove existence statements without actually exhibiting a particular value.

Ex: There exist irrational a and b with a^b rational.

Pf: Consider $(\sqrt{2})^{\sqrt{2}}$. If it is rational, we are done. If it is irrational, set $a = (\sqrt{2})^{\sqrt{2}}$ and $b = \sqrt{2}$, so that

$$((\sqrt{2})^{\sqrt{2}})^{\sqrt{2}} = (\sqrt{2})^{\sqrt{2} \cdot \sqrt{2}} = (\sqrt{2})^2 = 2. \quad \square$$

Buried in here is an idea that P is either true for x or it is false, and there is no third course. (This is different from demonstrating either of these, which is quite subtle.) This is usually summarized as saying that if P is ~~not~~ not-false, then it is true, and conversely.

This leads to a different kind of proof technique: contradiction. The idea is that if some premise leads you to say that something else must be both true and false, then your premise itself must have been unsound.

Ex. There are infinitely many prime numbers.

Pf. Suppose otherwise, that there are just finitely many, named p_1, p_2, \dots, p_k . We then form the number $N = (p_1 \cdot p_2 \cdot \dots \cdot p_k) + 1$, which is not divisible by p_j for any j .

This means that either N is prime (and not on the list) or that N decomposes into primes not on the list. In either case, we have shown our complete list of primes to be incomplete — a contradiction. Our initial assumption must have been wrong: it must instead be the case that there are ∞ many prime numbers. \square

Functions, properties, cardinalities

We have one more foundational issue to address before we begin linear algebra in earnest.

We will avoid actually saying what a set is. Suffice it to say that it is a collection of elements for which membership can be tested, e.g., $2 \in \{n \in \mathbb{N} \mid n \text{ is even}\} \subseteq \mathbb{N}$,
 $3 \notin$

A function $f: A \rightarrow B$ is an assignment of elements of A to those of B . That is, for any element $a \in A$ there is a single corresponding element $f(a) \in B$, called its image.

Functions tend to serve two purposes: operation and transmogrification.

Ex: The operation "+" on \mathbb{R} can be thought of as a function $+: \mathbb{R} \times \mathbb{R} \rightarrow \mathbb{R}$, where " $\mathbb{R} \times \mathbb{R}$ " indicates the set of pairs of real numbers. You can also specialize this to get a function $s(x) = x + 1$ by setting $y = 1$. There's also $p(x) = x^2$, which comes from specializing $+: \mathbb{R} \times \mathbb{R} \rightarrow \mathbb{R}$ to $x = y$.

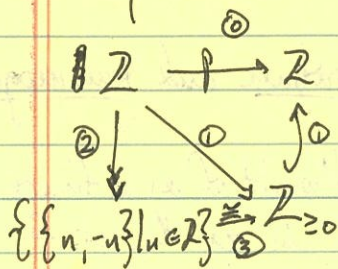
Ex: The ~~oper~~ function $\cos: \mathbb{R} \rightarrow \mathbb{R}$ plays more of the second role: it takes in an angle value (thought of as a real number) and gives out a ratio of lengths (thought of as a real number).

These functions ~~but~~ have certain properties which tell you interesting information about them.

• Injectivity: A function is injective if no two inputs give the same output. For transmogrification, this is a losslessness, that you can recover the input uniquely from the output. (\cos is not injective, because $\cos(0) = \cos(2\pi)$.) For operations, this is about solutions: s is injective, so $x + 1 = s(x) = y$ has a single solution. p is not, so $x^2 = p(x) = y$ may have many.

- Surjectivity: This is the statement that every output has at least one input realizing it. For transformations, this is about efficiency: there's no "wasted space" in the codomain of impossible values. For operations, this is about solutions: $x+1 = c(x) = y$ is surjective, so it is always possible to solve this eqⁿ for x , no matter what y is. $x^2 = p(x) = y$ is not, so there are y with no solutions in x .
- Bijectivity: Simultaneously injective and surjective. There are "perfect dictionary" transformations, or equations with exactly 1 solution for any choice of y .

Every function can be broken into these parts in the following way:



① If the function is not surjective, then we can restrict its codomain to just the values it does take on. In turn, this subset injects into the original codomain.

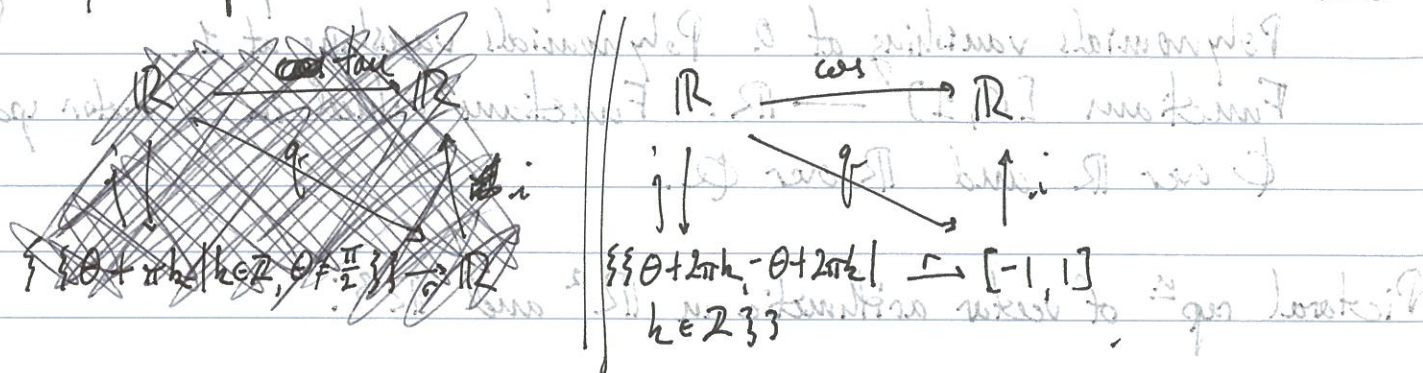
② If the function is not injective, we can collect together all the elements of the domain that give the same value into different subsets. These sets are said to partition the domain, meaning they do not overlap + yet their union is all of the domain. There is a surjective map assigning each element to the subset it belongs to.

③ Finally, the original function defines a new function on at the bottom: given a subset, the f^u takes on the same value on any of its members, so gives an element of the restricted codomain. This map is surjective and injective, hence bijective.

This is a good repⁿ of what functions "do". They forget a little information, ^{eg, the sign} then they represent what's left inside of the codomain according to some rule.

Ex: ~~function~~ $f(x) = \cos(x)$. ② tells you it's 2π -periodic + ③ tells you it lies in $[-1, 1]$.

First, a stray example from last time:



Vector Spaces

Remember that we are interested in functions $T: V \rightarrow W$ satisfying equations like $T(k \cdot u) = k \cdot T(u)$ and $T(u + v) = T(u) + T(v)$, where V and W are fancy w/domains. We need to make sense of "+" and "." inside of V and W .

~~Def~~ Ex: Represent a point in the plane by a coordinate pair (x, y) . Then rotation by 90° is specified by $T \begin{pmatrix} x \\ y \end{pmatrix} = \begin{pmatrix} -y \\ x \end{pmatrix}$. *same objects*

Defining + and \cdot componentwise $\begin{pmatrix} x \\ y \end{pmatrix} + \begin{pmatrix} w \\ z \end{pmatrix} = \begin{pmatrix} x+w \\ y+z \end{pmatrix}$, $k \begin{pmatrix} x \\ y \end{pmatrix} = \begin{pmatrix} kx \\ ky \end{pmatrix}$, *different!*

we find that T is a linear map. However, you can see that the op^{ns} + and \cdot are kind of complicated!

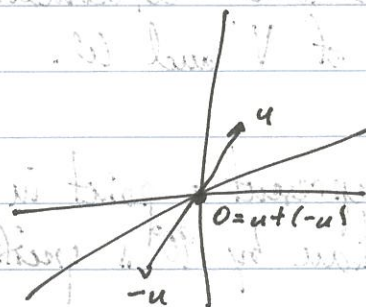
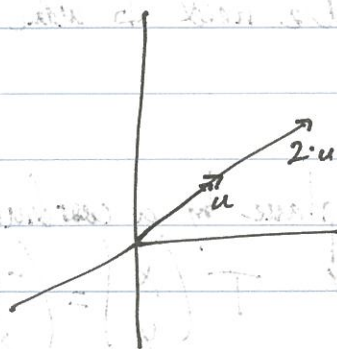
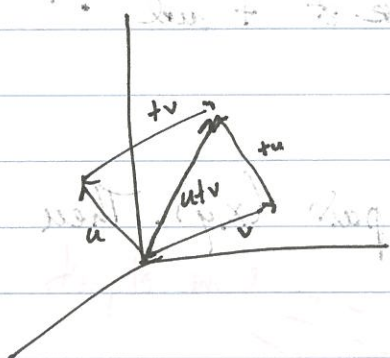
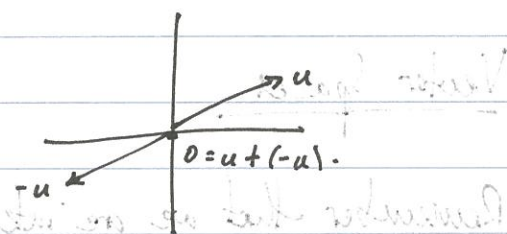
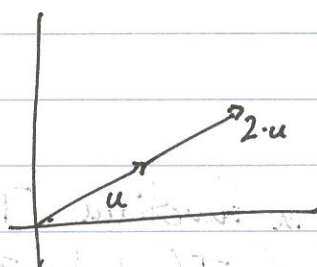
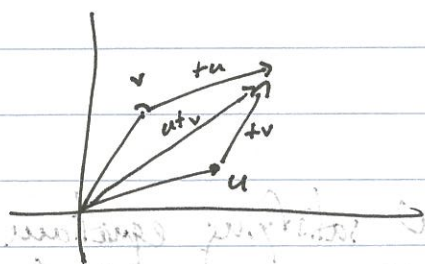
Def: A field k is a set with $+$, $-$, \cdot , and $/$ \Leftarrow defined on nonzero elements satisfying comm., assoc. + distributivity.

Ex: \mathbb{R} , \mathbb{C} , \mathbb{Z}/p , \mathbb{Q} . Non-ex: \mathbb{N} , \mathbb{Z} , $\mathbb{Z}/4$.

Def: A vector space V (over k) is a set with $+, -: V \times V \rightarrow V$ and $\cdot: k \times V \rightarrow V$, also satisfying comm. assoc. + distributivity

Ex: \mathbb{R}^n and \mathbb{C}^n . \mathbb{R}^∞ . ~~Polynomials~~ Polynomials. Polynomials of degree n .
 Polynomials vanishing at 0. Polynomials vanishing at 1.
 Functions $[0, 1] \rightarrow \mathbb{R}$. Functions valued in a vector space.
 \mathbb{C} over \mathbb{R} and \mathbb{R} over \mathbb{Q} .

Pictorial repⁿ of vector arithmetic in \mathbb{R}^2 and \mathbb{R}^3 :



Subspaces (1.c)

Implicit in our discussion thus far has been a notion of a subset:

a subset $Y \subseteq X$ is a set s.t. each element $y \in Y$ is already also an elt. $y \in X$.

This is a statement about size: X is at least as large as Y .

They're often described by properties: $\{n \in \mathbb{N} \mid n \text{ is div. by } 2\} \subseteq \mathbb{N}$,
the subset of even natural numbers.

There's a corresponding notion for vector spaces: $U \subseteq V$ is a subspace of V
if U is itself a vector space with the same op^{ns} as on V .
a subset and

Ex: $\{(x, y, z) \in \mathbb{R}^3 \mid x = y\} \subseteq \mathbb{R}^3$.

Polynomials vanishing at 0 \subseteq all polynomials

Non-ex: $\{(x, y, z) \in \mathbb{R}^3 \mid x = 5\} \subseteq \mathbb{R}^3$ is a subset but not a subspace.

This is because $(5, 0, 0)$ and $(5, 10, 32)$ are elements of the subset, but $(5, 0, 0) + (5, 10, 32) = (10, 10, 32)$ is not.

Sets have various interesting op^{ns} on them, like intersection, union, + complement.
These have analogues in vector spaces, but their behavior is more complex.

Intersection: The intersection of 2 subspaces is a subspace.

Union: The union of 2 subspaces U_1, U_2 is a subspace iff one contains the other. (This is homework) This has a replacement though: the sum is $U_1 + U_2 = \{u_1 + u_2 \mid u_1 \in U_1, u_2 \in U_2\}$.

Lemma: This is the smallest subspace containing U_1 and U_2 .

Pf:

Pf: It is a subspace: $(u_1 + u_2) + (u'_1 + u'_2) = \overbrace{(u_1 + u'_1)}^{eU_1} + \overbrace{(u_2 + u'_2)}^{eU_2}$, and
 $k \cdot (u_1 + u_2) = k \cdot u_1 + k \cdot u_2$. "Smallest" means that any other
subspace W with $U_1, U_2 \subseteq W$ has $U_1 + U_2 \subseteq W$. This is clear
too: for $u_1 + u_2 \in U_1 + U_2$, $u_1 \in U_1$ and $u_2 \in U_2$ and hence
 $u_1, u_2 \in W$. Then $u_1 + u_2 \in W$ b/c W is a subspace, and $U_1 + U_2 \subseteq W$. \square

Direct sum: A particularly nice kind of sum of subspaces is when $U_1 \cap U_2 = \{0\}$.

In this case, any $v \in U_1 + U_2$ has a unique representation as $v = u_1 + u_2$.

Pf: If $v = u_1 + u_2$ and $v = u'_1 + u'_2$, then $v - v = u_1 + u_2 - u'_1 - u'_2 = 0$,

and $\underbrace{u_1 - u'_1}_{eU_1, \neq 0} = \underbrace{u_2 - u'_2}_{eU_2, \neq 0}$.

This violates the intersection condition. \square

This is a kind of "disjoint union" condition: $U_1 + U_2$ have no overlap.

Complementation: For $A \subseteq X$ a subset, there is another ^{unique} subset $X \setminus A$
such that $A + X \setminus A$ are disjoint and $A \cup (X \setminus A) = X$.

This is kind of true for vector spaces — what fails is unicity.

Proving ~~the~~ existence in generality is more trouble than it's worth —
you need the Axiom of Choice. Instead, let's look at how unicity fails.

Ex: $U := \{(x, y) \in \mathbb{R}^2 \mid x = y\} \subseteq \mathbb{R}^2$.

One complement: $W = \{(x, y) \in \mathbb{R}^2 \mid x = -y\}$.

Pf: If $x = y$ and $x = -y$, then $x = y = 0$,

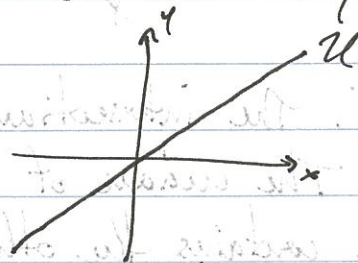
so $U \cap W = \{0\}$. Given $(s, t) \in \mathbb{R}^2$, we solve $x + \frac{x}{2} = s$, $x - \frac{x}{2} = t$

to get $x = \frac{s+t}{2}$, $x' = \frac{s-t}{2}$.

Ex': $W = \{(\frac{x}{2}, y) \in \mathbb{R}^2 \mid x = 0\}$. Pf: Again, $U \cap W = \{0\}$. For (s, t) ,

we find $\begin{pmatrix} s \\ t \end{pmatrix} = \begin{pmatrix} s \\ s \end{pmatrix} + \begin{pmatrix} 0 \\ t-s \end{pmatrix}$

\cap
 $U \quad W$



Finite-dimensional vector spaces (2.*)

Last time we skirted around the existence of ~~orthogonal~~ complements of vector subspaces. Here's a naive approach to constructing a complement which will be of interest to us.

Procedure: Start with a subspace $U \subseteq V$ and $W = 0$.

① If $U + W \neq V$, there is some missing vector $v \in V \setminus (U + W)$.

② Replace W by $W + \langle v \rangle$, where $\langle v \rangle = \{c \cdot v \mid c \in K\}$ is the smallest subspace containing v .

③ ~~Go~~ Go back to ①.

→ If not, then we're done: $U \cap W = 0$ and $U + W = V$.

WARNING: THIS MAY NOT TERMINATE IF V IS "TOO LARGE"!!

The subspace W we construct has a very particular form:

$$W = \langle v_1 \rangle + \langle v_2 \rangle + \dots + \langle v_n \rangle = \{c_1 v_1 + \dots + c_n v_n \mid c_i \in K\}$$

where v_i is the vector picked on the i^{th} time through the loop.

Def: W is called the span of (v_1, \dots, v_n) . A particular element $w = c_1 v_1 + \dots + c_n v_n$ is called a linear combination of (v_1, \dots, v_n) .

There is an interesting edge case of this algorithm: if $U = 0$, then its complement should be all of V . However, the algorithm presents V in a special form: $V = \langle v_1 \rangle + \langle v_2 \rangle + \dots + \langle v_n \rangle$. If the algorithm terminates, V is called finite dimensional (of dimension n).

Ex: $\mathbb{R}^2 \cong \langle (1, 0) \rangle + \langle (0, 1) \rangle$.

$\{\text{polynomials}\}$ is not finite dimensional. $(1, x, x^2, x^3, \dots)$

$\{\text{polynomials of degree} \leq n\} = \langle 1 \rangle + \langle x \rangle + \dots + \langle x^n \rangle$ has dimension $(n+1)$.

There is a flaw in using this algorithm as a definition: it is non-deterministic, meaning that it may behave differently based on what v_j is chosen at each step. This is worrying: does it sometimes terminate & sometimes not? Is the concluding number n always the same? We will have to work for a while to see this.

Lemma: A linear dependence is a nonzero linear combination $c_1 w_1 + \dots + c_d w_d = 0$.

Suppose c_j is nonzero. Then $\text{span}(w_1, \dots, w_d) = \text{span}(w_1, \dots, \overset{A}{w_{j-1}}, \overset{B}{w_{j+1}}, \dots, w_d)$, with w_j removed.

Pf: Automatically, $B \subseteq A$. To see $A \subseteq B$, note that

$w_j = \frac{1}{c_j} (-c_1 w_1 - \dots - c_{j-1} w_{j-1} - c_{j+1} w_{j+1} - \dots - c_d w_d)$, which lets us write any element $a \in A$ as

$$a = k_1 w_1 + \dots + k_d w_d$$

$$= k_1 w_1 + \dots + k_{j-1} w_{j-1} + k_j (\text{---}) + k_{j+1} w_{j+1} + \dots + k_d w_d.$$

This does not involve w_j . \square

Rem: Being linearly independent is the same as $\langle w_1 \rangle + \dots + \langle w_d \rangle$ being a direct sum.

Cor: The length of any linearly independent list \leq the length of any spanning list. $\leftarrow v_s, w_s$

Pf: ① Start with (w_1, \dots, w_d) and (v_1, \dots, v_n) .

② Prepend the first v -vector to the w list.

③ There is a dependence, ^{in the w -list} not involving the v 's. Use this to eliminate a w -vector. (\uparrow any finite set of them is still lin. ind.)

④ Repeat.

Eventually, you'll run out of v 's, before you run out of w 's.

That means $n \leq d$. \square

Cor: The algorithm gives the same n no matter what.

Pf: If you have v_1, \dots, v_n and $v'_1, \dots, v'_{n'}$, then $n \leq n'$ and $n' \leq n$. \square

More on the Dimensional algorithm: (2.)*

We can squeeze some more out of the ideas from last time:

Cor. If $U \subseteq V$ is a subspace + V is finite dim., then so is U .

Pf. Run the algorithm on U and on V . The list resulting from U is linearly independent (in V) and the list from V spans V .

The Lemma from last time says $\text{length span} \geq \text{length l.i.}$ \square
(In fact $\dim U \leq \dim V$.)

I've been obtuse and avoided giving you some useful vocabulary:
a basis for V is a set that is both linearly independent and spans V . (The list resulting from the algorithm is a basis for the complementary subspace.)

Reinforcing ex. $\{(1,0), (0,1)\}$ is a basis for \mathbb{R}^2 . None of $\{(3,5), (2,1)\}$. However, ~~neither~~ $\{(1,1)\}$ doesn't span
 $\{(1,0), (0,1), (1,1)\}$, ~~and~~ $\{(1,0), (2,0)\}$ are bases. linearly dependent!

Lemma: Any spanning list can be shortened to a basis.

Pf. ① Start with $j=1$.

② If v_j is in $\text{span } \{v_1, \dots, v_{j-1}\}$, then discard v_j .
Otherwise keep it and continue to the next j .

At the end, the rest of the list will still span V . It's now linearly independent: if there were a dependence, then there would be a last nonzero coefficient in the dependence. That would violate step ② at that stage. \square

~~Cor:~~

Lem: Every l.i. list of vectors in a finite dim^t V extends to a basis of V .

Pf: Take $U = \text{span}\{u_1, \dots, u_d\}$ to be the span of the l.i. list. Use the complementary algorithm to find $W = \langle v_1 \rangle + \dots + \langle v_n \rangle$. Then $V = U + W = \langle u_1 \rangle + \dots + \langle u_d \rangle + \langle v_1 \rangle + \dots + \langle v_n \rangle$, and this is a direct sum because the lists are l.i. and $U \cap W = 0$. \square

~~Lem:~~

Lem: If $\{v_1, \dots, v_n\}$ is l.i. and $\dim V = n$, then $\{v_i\}$ is a basis.

Pf: Extend it to a basis — but it's already length n ! So no new vectors are added. \square

Lem: If $\{w_1, \dots, w_d\}$ is spanning and $\dim V = d$, then $\{w_i\}$ is a basis.

Pf: It can be reduced to a basis — but it's already length d ! So, no vectors can be erased. \square

← Exer 2.43

Lem: For $U_1, U_2 \subseteq U$, we have $\dim(U_1 + U_2) = \dim U_1 + \dim U_2 - \dim(U_1 \cap U_2)$.

Pf: Use the algorithm to build a basis ^{5.4.3} for $U_1 \cap U_2$. Consider it as a l.i. set in U_1 and U_2 separately, and extend it to a basis $\{u_1, \dots, u_d, v_1, \dots, v_n\}$ of U_1 and $\{u_1, \dots, u_d, w_1, \dots, w_m\}$ of U_2 . We claim that the combined list $\underbrace{\{u_1, \dots, u_d\}}_{U_1 \cap U_2} \underbrace{\{v_1, \dots, v_n\}}_{\text{new to } U_1} \underbrace{\{w_1, \dots, w_m\}}_{\text{new to } U_2}$ is a basis for $U_1 + U_2$.

It clearly spans. Suppose there were a linear dependence:

$$a_1 u_1 + \dots + a_d u_d + b_1 v_1 + \dots + b_n v_n + c_1 w_1 + \dots + c_m w_m = 0.$$

$$-(a_1 u_1 + \dots + a_d u_d + b_1 v_1 + \dots + b_n v_n) = c_1 w_1 + \dots + c_m w_m.$$

$\in U_1 \cap U_2$

$\in U_1$

$\in U_2$

← actually $U_1 \cap U_2$.

Hence $c_1 w_1 + \dots + c_m w_m = d_1 u_1 + \dots + d_d u_d$. Substituting this back:

$$a'_1 u_1 + \dots + a'_d u_d + b_1 v_1 + \dots + b_n v_n = 0, \text{ but this list is l.i.} \quad \square$$

~~Linear maps & Factorization~~

Linear maps + Kernels (3.A-B)

Finally, we turn our attention to how vector spaces relate to one another through linear maps. One more time:

Defⁿ: A fⁿ $T: V \rightarrow W$ is linear (T, W vector spaces) when $T(v_1 + v_2) = T(v_1) + T(v_2)$ and $T(k \cdot v) = k \cdot T(v)$.

Ex: ① $T: \mathbb{R}^2 \rightarrow \mathbb{R}$ given by $T(x, y) = x - y$, or
 $T': \mathbb{R}^2 \rightarrow \mathbb{R}$ given by $T(x, y) = y$.

② $T: \{\text{polynomials}\} \rightarrow \mathbb{R}$ given by $T(f) = f(0)$, or
 $T': \{\text{polynomials}\} \rightarrow \mathbb{R}$ given by $T(f) = f(1)$.

The basic op^s on linear functions are:

① Addition: given $T_1, T_2: V \rightarrow W$, we can form $(T_1 + T_2)(v) = T_1(v) + T_2(v)$, which is also linear.

② Scaling: given $T: V \rightarrow W$ and $k \in K$, we can form $(k \cdot T)(v) = k \cdot (T(v))$, which is also linear.

③ Composition: given $V \xrightarrow{T} W \xrightarrow{T'} W'$, we can compose $(T' \circ T)(v) = T'(T(v)) \in W'$ to get a linear map.

These play nicely with each other. For instance, \circ distributes over $+$.

There ~~are~~ also 2 natural subspaces associated to T :

Def: The kernel of $T: V \rightarrow W$ is $\{v \in V \mid T(v) = 0\} \subseteq V$.

It is a subspace. The image of T is $\{w \in W \mid \exists v \in V \text{ with } T(v) = w\}$

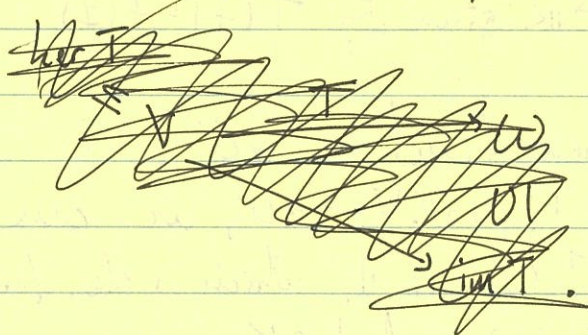
It, too, is a subspace.

\cap
 W

<p>Ex: ① $\ker T = \{(x, y) \in \mathbb{R}^2 \mid x - y = 0\} \subseteq \mathbb{R}^2$, $\ker T' = \{(x, y) \in \mathbb{R}^2 \mid y = 0\} \subseteq \mathbb{R}^2$. ② $\ker T = \{f \text{ a poly}^{\mathbb{C}} \mid f(0) = 0\} \subseteq \{\text{all polynomials}\}$, $\ker T' = \{f \text{ a poly}^{\mathbb{C}} \mid f(1) = 0\} \subseteq \{\text{all polynomials}\}$.</p>	<table border="1" style="width: 100%; border-collapse: collapse; text-align: center;"> <tr> <td style="padding: 5px;">$\dim \ker$</td> <td style="padding: 5px;">$\dim V$</td> <td style="padding: 5px;">$\dim \text{im}$</td> </tr> <tr> <td style="padding: 5px;">1</td> <td style="padding: 5px;">2</td> <td style="padding: 5px;">1</td> </tr> <tr> <td style="padding: 5px;">line</td> <td style="padding: 5px;">\mathbb{R}^2</td> <td style="padding: 5px;">line</td> </tr> <tr> <td style="padding: 5px;">line</td> <td style="padding: 5px;">\mathbb{R}^2</td> <td style="padding: 5px;">line</td> </tr> </table>	$\dim \ker$	$\dim V$	$\dim \text{im}$	1	2	1	line	\mathbb{R}^2	line	line	\mathbb{R}^2	line
$\dim \ker$	$\dim V$	$\dim \text{im}$											
1	2	1											
line	\mathbb{R}^2	line											
line	\mathbb{R}^2	line											

These are all subspaces we've thought about before! This is interesting: what exactly is the relationship between $f \stackrel{u}{=} T: V \rightarrow W$ and subspaces U ? Can we get all U ? How many T s give the same U ? What does W have to do with it? (Consider $W=0$.)

~~As the moment, we're going to think about something we've seen before: factorizations. The image subspace plays the role of "Step 1".~~



These are all interesting questions. For the moment, we're going to produce a relation between $\ker T$ and $\text{im } T$:

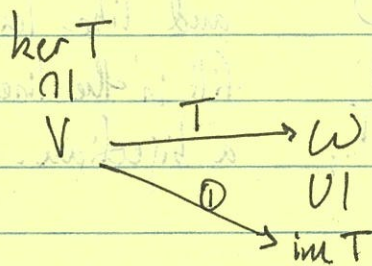
Lemma: $\dim(\ker T) + \dim(\text{im } T) = \dim V$ for f.d. V .

Pf: Extend a basis of $\ker T$ to one of V . The image of the extension in $\text{im } T$ is a basis there. \square

This is interesting: the proof says that $\text{im } V$ has the same dimension as a complement of $\ker T$ — but $V \rightarrow \text{im } T \hookrightarrow W$ is canonical, whereas a choice of complement $(\ker T)^c \leq V$ is not unique. We will think about this next time in the context of factorizations.

Factorizations for Linear Maps:

Last time we talked about subspaces associated to



We basically did Step ① by factoring T through $\text{im } T$, which injects into W .

~~We're missing an analogue of Step ②: a way to build a surjection with kernel $\ker T$.~~

To start, this picture suggests an interesting lemma:

Lem: A map $T: V \rightarrow W$ is injective if and only if $\ker T = 0$.

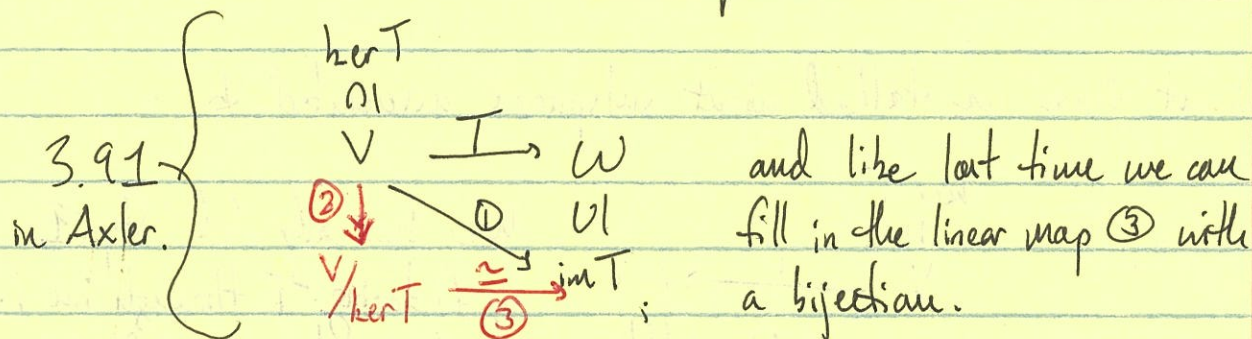
Pf: If T is injective, then $\{v \in V \mid T(v) = 0\} = 0$. If T is not injective, then there are $v_1 \neq v_2$ with $T(v_1) = T(v_2)$. But then $T(v_1 - v_2) = T(v_1) - T(v_2) = 0$ exhibits $v_1 - v_2 \in \ker T$. \square

To complete the picture, we're missing an analogue of Step ②: a way to build a surjection with kernel the subspace $\ker T \leq V$. This actually looks a lot like what we did for sets:

Defⁿ: Given $U \leq V$, we define V/U by $V/U = \{v + U \mid v \in V\}$, a collection of subsets of V .

Lem: There is a map $f: V \rightarrow V/U$ given by $f(v) = v + U$ which is surjective with kernel U . \square

This construction fills in the 2nd step:



Let's think about the members of V/U some more.

Rem: U itself is one member, since $0+U=U$.

Rem: The other members of V/U look like translates of U off of the origin. We know there are not subspaces, but they are useful enough to earn a name: they are affine subspaces (or translates).

these show up when considering the solⁿ set to equations like $T(v)=w$.

Lem: The solⁿ set $\{v \in V \mid T(v)=w\}$ is empty or a translate of $\ker T$.

Pf: If the solⁿ set is empty, we are done. If it's nonempty, pick a $v \in V$ with $T(v)=w$. Then $v + \ker T$ is exactly the solⁿ set:

① For $h \in \ker T$, $T(v+h) = T(v) + T(h) = w + 0 = w$ gives another solⁿ.

② For another solⁿ v' , $T(v-v') = T(v) - T(v') = w - w = 0$, so $v-v' \in \ker T$. \square

Ex: y -axis
 $\mathbb{R}^2 \xrightarrow{\text{proj}_x} \mathbb{R}^2$
 \downarrow
 space of vert. lines $\xrightarrow{x\text{-coord}}$ $\mathbb{R} = \{x, 0 \mid x \in \mathbb{R}\}$.

Degeneracy: If $\ker T = 0$, then $V \rightarrow V/\ker T$ is already bijective. If T is surjective, then $\text{im } T = W$. Now take $\dim V = \dim W < \infty$.

• If T is bijective, then $T \text{ surj} \Rightarrow \dim \ker T = 0 \Rightarrow T \text{ inj.}$

• If T is not bijective, then $T \text{ inj} \Rightarrow$ a basis of $\text{im } T$ gives a l.i. set in W of size $= \dim W$.

Bases as presentations

A few times in this class we've drawn some picture like

$$\begin{array}{ccc} V & \xrightarrow{f} & W \\ \downarrow j & \xrightarrow{i \circ \tilde{f} \circ j} & \uparrow i \\ V/\ker f & \xrightarrow{\tilde{f}} & \text{im } f \end{array} \quad \text{to communicate the identity } f = i \circ \tilde{f} \circ j.$$

These pictures are called diagrams, their nodes are labeled by vector spaces, their arrows by linear maps, and they encode how different paths with the same start + end are the same.

A useful puzzle piece when drawing these pictures is the isomorphism, which is a bijection (or invertible) linear map.

These look like

$$\begin{array}{ccc} V & \xrightarrow{f} & W \\ \xleftarrow{f^{-1}} & & \end{array}$$

Note that going $V \xrightarrow{f} W \xrightarrow{f^{-1}} V$ is the same as staying stationary at $V \xrightarrow{\text{i.e., the identity map } 1: V \rightarrow V}$.

Lemma: A basis S for a vector space V gives an isom $K^n \xrightarrow{\varphi} V$, ~~and~~ and conversely.

Pf: Given $S = (v_1, \dots, v_n)$, we define $\varphi(k_1, \dots, k_n) = k_1 v_1 + \dots + k_n v_n$.

This is surjective because S spans, and it's injective because S is linearly independent. If we're instead given φ , we set $v_j = \varphi(\underbrace{0, \dots, 0, 1, 0, \dots, 0}_{e_j})$ in the j^{th} position. Again, surjectivity gives spanning + injectivity gives linear independence. \square

e_j

This means that K^n are the "standard" vector spaces, — and because there are such concrete spaces, we can say a lot about them.

~~Lemma: Linear maps $V \rightarrow W$ are determined by their values on a basis.~~
~~Lemma: Linear maps $K^n \rightarrow K^m$ are encoded by $m \times n$ matrices.~~

~~Pf:~~

Lemma: Linear maps $K^n \rightarrow K^m$ are encoded by $m \times n$ matrices.

Pf: Any vector $v = (k_1, \dots, k_n) \in K^n$ can be decomposed as

$v = k_1 e_1 + \dots + k_n e_n$, where e_j is as before. So, we only need to evaluate $f(e_j)$, which itself has a decomposition

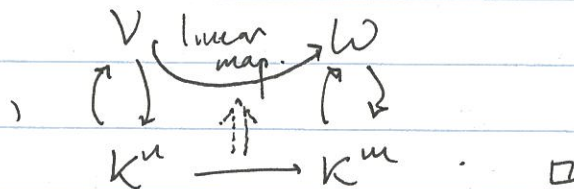
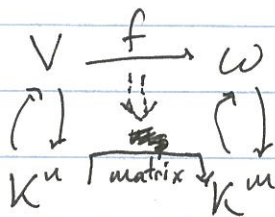
$f(e_j) = a_{1j} f_1 + a_{2j} f_2 + \dots + a_{mj} f_m$, f_j the j -th standard basis vector in K^m .

Arranging these numbers into a grid (a_{ij}) , we uncover a matrix. ~~Conversely~~, linear indep. of (e_j) shows that ~~this~~ a matrix specifies a function, which is checked to be linear.

Linear indep. of (f_j) shows that no two matrices rep the same map. \square

Theorem: Under a choice of basis in the domain and codomain, linear maps and matrices correspond.

Pf:



Lemma: Matrix multiplication encodes function composition.

Pf: Represent f by (a_{ij}) and g by (b_{kl}) . Then

$$\begin{aligned} f(g(e_l)) &= f\left(\sum_{k=1}^m b_{kl} f_k\right) = \sum_{k=1}^m b_{kl} f(f_k) = \sum_{k=1}^m b_{kl} \left(\sum_{i=1}^m a_{ik} g_i\right) \\ &= \sum_{i=1}^m \left(\sum_{k=1}^m a_{ik} b_{kl}\right) g_i = (c_{il}). \quad \square \end{aligned}$$

Duality (J.F)

One of the consequences of last time is:

Cor: $\dim L(V, W) = \dim V \cdot \dim W$.

"the dual"

In particular, $\dim L(V, K) \cong \dim V$, yet $V^* := L(V, K)$ has interesting properties not exactly like those of V . The most basic such property is that it is backwards or contravariant.

Def: A map $f: V \rightarrow W$ induces a map $f^*: W^* \rightarrow V^*$ defined by $(f^*\varphi)(v) = \varphi(f(v))$ (or precomposition).

Ex: There is an iso $L(K^n, K) \cong K^n$ by $\varphi \mapsto (\varphi(e_j))_j$.

The induced matrix is $(a_{ij})^* = (a_{ji})$, called the transpose.

You might enjoy checking the identity $(AB)^* = B^*A^*$.

We have two tasks to take care of today.

I. Pairing: A map $V \times W \xrightarrow{\langle \cdot, \cdot \rangle} K$ is called a pairing, and the pairing is moreover perfect if $\forall w \in W \exists v \in V$ with $\langle v, w \rangle \neq 0$.

Lemma: A perfect pairing determines an injection $W \hookrightarrow V^*$.

Pf: This is just "currying": $\tau(w)(v) = \langle v, w \rangle$, and the perfection condition shows that $\tau(w) \neq 0$ so that τ is injective. \square

~~Cor~~: Lemma: There is a "natural" iso $V \cong V^{**}$. There is an evaluation perfect pairing $V \times V^* \rightarrow K$, and flipping this around gives an injection $V \hookrightarrow (V^*)^*$. Since these are equidimensional, it's an isomorphism. (In the ∞ -dimensional case, we at least get an injection.)

II. Subspaces associated to dual maps

Continuing our obsession with subspaces, it would be nice to understand $\ker(f^*)$ and $\text{im}(f^*)$ in terms of f .

Toward this, we make the following interrelating def^u:

Def: Given $U \subseteq V$, we define the annihilator $U^\circ \subseteq V^*$ by $\{\varphi \in V^* \mid U \subseteq \ker(\varphi), \text{ or } \varphi(U) = 0\}$. This is a subspace.

Lemma: $\dim U + \dim U^\circ = \dim V$.

Pf: Consider $i: U \hookrightarrow V$ and its dual $i^*: V^* \longrightarrow U^*$.

We have $\dim V^* = \dim \ker(i^*) + \dim \operatorname{im}(i^*)$

$$\dim V^* = \dim U^\circ + \dim U^*$$

$$\dim V = \dim U^\circ + \dim U. \quad \square$$

The annihilator also gives the desired relations between f and f^* :

Lemma: $\ker(f^*) = (\operatorname{im} f)^\circ$, and $\operatorname{im}(f^*) = (\ker f)^\circ$.

Pf: The first equality is a matter of definitions. In the second case, only $\operatorname{im}(f^*) \subseteq (\ker f)^\circ$ is obvious from the def^u. However, $\dim \operatorname{im} f^* = \dim W^* - \dim \ker f^*$

$$= \dim W - \dim (\operatorname{im} f)^\circ = \dim \operatorname{im} f$$

$$= \dim V - \dim \ker f = \dim (\ker f)^\circ. \text{ So, } \operatorname{im}(f^*)$$

is a top-dim^d subspace of $(\ker f)^\circ$ and hence they're equal. \square

Cor: $\dim \operatorname{im} f^* = \dim W^* - \dim \ker f^*$

$$= \dim W^* - \dim (\operatorname{im} f)^\circ$$

$$= \dim \operatorname{im} f. \quad \square$$

Polynomials over \mathbb{R} and \mathbb{C} (Ch. 4)

Soon, we will move on to the second major goal of this course: understanding eq⁴ of the form $f(v) = k \cdot v$ for a linear f^n $f: V \rightarrow V$ from a vector space to itself. We will find out that analysis of this situation ~~pro~~ involves polynomials, which have particularly nice properties / \mathbb{C} + \mathbb{R} .

Lemma: The only zero function is the zero polynomial (over \mathbb{R} or \mathbb{C}).

Pf: Suppose $f(x)$ is a nonzero polynomial $f(x) = a_m x^m + \dots + a_1 x + a_0$.

May as well take $a_m = 1$, and set $z = |a_0| + |a_1| + \dots + |a_{m-1}| + 1$.

We must have $|z| > 1$, so $z^{m-1} \leq z^{m-1}$, so

$$|a_0 + a_1 z + \dots + a_{m-1} z^{m-1}| \leq |a_0| + |a_1| z + \dots + |a_{m-1}| z^{m-1}$$

$$= (|a_0| + |a_1| + \dots + |a_{m-1}|) z^{m-1}$$

$$< (|a_0| + |a_1| + \dots + |a_{m-1}| + 1) z^{m-1} = z^m$$

Hence, $(a_0 + a_1 z + \dots + a_{m-1} z^{m-1}) + z^m \neq 0$. \square

Lemma: For $p, s \in \mathbb{P}$ ~~\mathbb{R}~~ ^{unique!} with $s \neq 0$ there are polynomials $q, r \in \mathbb{P}$ ~~\mathbb{R}~~ with $p = sq + r$ and $\deg r < \deg s$.

Pf: $T(q, r): \mathbb{P}_{n-m} \times \mathbb{P}_{m-1} \rightarrow \mathbb{P}_n$

$\underbrace{\quad}_{n} \underbrace{\quad}_{m-1} (q, r) \mapsto sq + r$ is a linear map.

For $\deg p \geq \deg s = m$, T is injective, since otherwise $sq = -r$ for nonzero poly^s of degrees $\geq m$ and $\leq m-1$.

But $\dim(\mathbb{P}_{n-m} \times \mathbb{P}_{m-1}) = (n-m+1) + (m-1+1) = n+1 = \dim \mathbb{P}_n$, so it is also ~~surjective~~ surjective. \square

A special case of this is when $\deg s = 1$.

Defⁿ: A root of a polynomial p is a value α with $p(\alpha) = 0$.

Cor. α is a root of p iff $(z - \alpha)$ divides $p(z)$.

Pf: If $p(z) = (z - \alpha)q(z)$, then $p(\alpha) = 0$. Otherwise, $p(z) = (z - \alpha)q(z) + r$ for some r with $\deg r = 0$, i.e., a constant. So, $p(\alpha) = r \neq 0$ and α is not a root. \square

Cor. A ^{nonzero} polynomial of $\deg n$ has at most n roots.

Pf: In degree zero, $f(z) = a_0 \neq 0$ has no roots. In degree 1, $f(z) = a_0 + a_1 z$ has a unique root $z = -a_0/a_1$. Otherwise, induct: in degree n $f(z)$ either has no roots (\checkmark) or at least one root. Pick one α and divide it out: $f(z) = q(z)(z - \alpha)$. The zero-product property reduces to q , with $\deg q = n - 1$. \square

Important fact: Every $\deg \geq 1$ poly^d over \mathbb{C} has a root. "Fund. Th^m of Algebra"

Cor: Every $f \in P(\mathbb{C})$ has a unique (up to order) factorization as $f(z) = c \cdot (z - \alpha_1) \cdots (z - \alpha_n)$.

Pf: The existence of fact^z follows from the Fact. If we had two such, we could pair one root by ZPP. The two quotients agree except maybe at α — but they must agree here too by the previous Cor. Induct. \square

Leun: Real polynomials factor into $(z - \alpha)$ and $(z - h)^2 + \beta^2$, $\beta > 0$.

Pf: Real roots occur in conjugate pairs: $p(\bar{\alpha}) = \overline{p(\alpha)} = \overline{0} = 0$. For a complex root, translate it to the origin to get $(z - (h + i\beta))(z - (h - i\beta))$. \square

Invariant Subspaces (S.A)

When we talked about matrices, we noticed that it was easier to understand $K^n \xrightarrow{f} K^m$ where $K_j = \langle v_j \rangle$ was the span of v_j .
 $K_1 \oplus \dots \oplus K_n \xrightarrow{f|_{K_1} \oplus \dots \oplus f|_{K_n}} K^m$

This idea holds in more generality: we can let the domain be V and take any sum decomposition of V . The idea, again, is that understanding $f|_{v_j}$ should be an easier problem than understanding f itself. Things are complicated by studying maps of the form $f: V \rightarrow V$ with the same domain + codomain.

Consider: $\mathbb{R}^2 \xrightarrow{\begin{pmatrix} 1 & 3 \\ 0 & 2 \end{pmatrix}} \mathbb{R}^2$
 $\downarrow \text{S.I.} \quad \downarrow \text{S.I.}$
 $\langle \begin{pmatrix} 1 \\ 0 \end{pmatrix} \rangle \oplus \langle \begin{pmatrix} 0 \\ 1 \end{pmatrix} \rangle \xrightarrow{\begin{pmatrix} 1 & 3 \\ 0 & 2 \end{pmatrix}} \langle \begin{pmatrix} 1 \\ 0 \end{pmatrix} \rangle \oplus \langle \begin{pmatrix} 0 \\ 1 \end{pmatrix} \rangle$
 The linear map in this example does not respect the subspace decomposition. The second factor of the domain is blurred across both in the source.

Defⁿ: A subspace $U \subseteq V$ is invariant if $f(U) \subseteq U$.

~~General prob~~

General problem: How can we find invariant subspaces? How finely can we find them (to avoid the trivial solⁿ $U=V$)?

We'll start as finely as possible.

Defⁿ: A vector $v \in V$ satisfying $f(v) = k \cdot v$ (i.e., $f(\langle v \rangle) \subseteq \langle v \rangle$) is called an eigenvector of eigenvalue k .

Ex: In the example above, $\begin{pmatrix} 1 \\ 0 \end{pmatrix}$ and $\begin{pmatrix} 3 \\ 1 \end{pmatrix}$ are eigenvectors of eigenvalues 1 and 2 respectively. Moreover, $\mathbb{R}^2 \cong \langle \begin{pmatrix} 1 \\ 0 \end{pmatrix} \rangle \oplus \langle \begin{pmatrix} 3 \\ 1 \end{pmatrix} \rangle$, so this is as fine as possible.

$$\begin{array}{ccc} \mathbb{R}^2 & \xrightarrow{\text{S.I.}} & \mathbb{R}^2 \\ \begin{pmatrix} 1 & 0 \\ 0 & 1 \end{pmatrix} & \xrightarrow{\begin{pmatrix} 1 & 0 \\ 0 & 2 \end{pmatrix}} & \begin{pmatrix} 1 & 0 \\ 0 & 1 \end{pmatrix} \\ \oplus & & \oplus \\ \begin{pmatrix} 3 & 1 \\ 0 & 1 \end{pmatrix} & \xrightarrow{\begin{pmatrix} 1 & 0 \\ 0 & 2 \end{pmatrix}} & \begin{pmatrix} 3 & 1 \\ 0 & 1 \end{pmatrix} \end{array}$$

Ex: Recall the rotation operator $\mathbb{R}^2 \xrightarrow{\begin{pmatrix} 0 & -1 \\ 1 & 0 \end{pmatrix}} \mathbb{R}^2$. This has no eigenvectors over \mathbb{R} .

Over \mathbb{C} , it does: $\begin{pmatrix} y \\ -x \end{pmatrix} = \begin{pmatrix} 0 & -1 \\ 1 & 0 \end{pmatrix} \begin{pmatrix} x \\ y \end{pmatrix} = \lambda \begin{pmatrix} x \\ y \end{pmatrix} = \begin{pmatrix} \lambda x \\ \lambda y \end{pmatrix}$ has nonzero solutions by setting $y = \lambda x$ and $-x = \lambda y = \lambda(\lambda x)$, or $\lambda^2 + 1 = 0$, or $\lambda = \pm i$.
Some corresponding eigenvectors are $(x, -ix)$ for i and (x, ix) for $-i$.

Lemma: If v_1, \dots, v_n are eigenvectors for distinct eigenvalues then they are l.i.

Pf: Let $a_1 v_1 + \dots + a_n v_n = 0$ be the earliest dependence. Then

$$\xrightarrow{f} a_1 \lambda_1 v_1 + \dots + a_n \lambda_n v_n = 0, \text{ and}$$

$$\xrightarrow{\cdot \lambda_k} a_1 \lambda_k v_1 + \dots + a_n \lambda_k v_n = 0, \text{ whose difference is an earlier dependence. } \square$$

Cor: $f: V \rightarrow V$ has at most $\dim V$ distinct eigenvalues. \square

Defⁿ: If $U \leq V$ is invariant for f , then we can build two operators:

$$\begin{array}{ccccc} U & \hookrightarrow & V & \longrightarrow & V/U \\ \text{the behavior of } f \text{ on } U & \rightsquigarrow & \downarrow f|_U & \downarrow f & \downarrow f/U \end{array} \quad \leftarrow \text{intuitively, what's left over ignoring } U.$$

Warning: Ignoring U can get you into trouble, if there's no invariant complement to U . Typical example:

$$\begin{array}{ccccc} \begin{pmatrix} 0 \\ 1 \end{pmatrix} & \hookrightarrow & \begin{pmatrix} 0 \\ 1 \end{pmatrix} & \hookrightarrow & x \\ \begin{pmatrix} 0 \\ 1 \end{pmatrix} & \hookrightarrow & \mathbb{R}^2 & \longrightarrow & \mathbb{R} \\ \downarrow 0 & & \downarrow \begin{pmatrix} 0 & 0 \\ 1 & 0 \end{pmatrix} & & \downarrow 0 \end{array} \quad \begin{array}{l} \text{has zero } f|_U \text{ and } f/U, \\ \text{but } f \text{ is nonzero.} \end{array}$$

Eigenvectors + U-T matrices (S.B)

Q: How are we supposed to find eigenvectors? How do we even know that they exist?

Thm: $f: V \rightarrow V$ for V ^{an n-dim} vector space / \mathbb{F} has an eigenvector.

Pf: Pick $v \neq 0$ and consider $\{v, fv, f^2v, \dots, f^n v\}$. They must have a dependence: $a_0 v + a_1 fv + \dots + a_n f^n v = 0$, for a_i nonzero coeff^s a_i .

The polynomial $p(f) = a_0 + a_1 f + \dots + a_n f^n$ factors as $p(f) = c(f - \lambda_1) \dots (f - \lambda_n)$, and we substitute this in: $c(f - \lambda_1) \dots (f - \lambda_n)v = 0$. One of these maps $f - \lambda_j$ must fail to be injective, i.e., $\exists w$ with $fw = \lambda_j w$. \square

Some days ago, we worked through the example $\begin{pmatrix} 1 & 3 \\ 0 & 2 \end{pmatrix}$, which had eigenvalues 1 and 2. This behavior is actually 'generic': the eigenvalues of an upper-triangular matrix lie on its diagonal. Additionally, every matrix admits (/ \mathbb{F}) an upper-triangular presentation.

Rem: In terms of invariant subspaces, the matrix of f in a basis (v_1, \dots, v_n) is upper triangular when $\text{span}(v_1, \dots, v_j)$ is invariant for each j .

Cor: Over \mathbb{F} , every operator f admits an upper-triangular presentation.

Pf: We will induct on n , as the result is trivially true at $n=1$. By the Theorem, let λ be an eigenvalue for f , and set $U = \text{im}(f - \lambda)$.

This is a proper subspace which is invariant under f : for $u \in U$, $f(u) = (f - \lambda)u + \lambda u$ decomposes as two things in U . Hence, we can find an upper-triangular basis for $f|_U$, (u_1, \dots, u_m) , which we extend to a basis $(u_1, \dots, u_m, v_1, \dots, v_n)$ of V . By hypothesis, $u_j \in \text{span}(u_1, \dots, u_j)$. For v_j , $f(v_j) = (f - \lambda)(v_j) + \lambda v_j \in \text{span}(U) + \text{span}(v_j) \subseteq \text{span}(u_1, \dots, u_m, v_1, \dots, v_j)$. \square

(Alternatively: An eigenvector $v_i \in V$ gives a matrix $\begin{pmatrix} \lambda_i & \text{stuff from lifting} \\ 0 & \text{study } f(v_i) \text{ and induct.} \\ \vdots & \\ 0 & \text{This is } \nabla. \end{pmatrix}$.)

Cor: An upper-triangular matrix is ~~non~~^{invertible} iff its diagonal entries are nonzero.

Pf: If λ_j are all nonzero, we can back-substitute to get $v_j \in \text{im } f$ for all v_j in the basis. But then $\dim \text{im } f = \dim V$. Conversely, if $\lambda_j = 0$ for some j , then $\text{im } f|_{v_1, \dots, v_j} \subseteq \text{span}(v_1, \dots, v_{j-1})$. This forces f not to be injective. \square

Cor: The eigenvalues of an upper-triangular matrix appear in its diagonal.

~~Cor~~ Pf: $\begin{pmatrix} \lambda_1 - \lambda & * & \dots & * \\ 0 & \ddots & & \vdots \\ \vdots & & \ddots & * \\ 0 & \dots & 0 & \lambda_n - \lambda \end{pmatrix}$ is non-invertible if and only if $\lambda = \lambda_j$ for some j . \square

Ex: $M = \begin{pmatrix} -2 & 3 \\ -4 & 5 \end{pmatrix}$. $v = \begin{pmatrix} 1 \\ 0 \end{pmatrix}$, $Mv = \begin{pmatrix} -2 \\ -4 \end{pmatrix}$, $M^2 v = \begin{pmatrix} -8 \\ -12 \end{pmatrix} = 3Mv - 2v$. $M^2 - 3M + 2 = 0$. $(M-1)(M-2) = 0$.

~~Then~~ $\lambda = 1$, as a guess. Has $v_1 = \begin{pmatrix} 1 \\ 0 \end{pmatrix}$ as a witness, so it's an eigenvalue.

Then $M - \lambda I = \begin{pmatrix} -3 & 3 \\ -4 & 4 \end{pmatrix}$ has image $U = \text{span}\left\{\begin{pmatrix} 1 \\ 1 \end{pmatrix}\right\} = U_1$.

Extend this to a basis $\left\{\begin{pmatrix} 1 \\ 1 \end{pmatrix} = u_1, \begin{pmatrix} 0 \\ 1 \end{pmatrix} = v_1\right\}$. Then

$$\begin{array}{ccc} \mathbb{R}^2 & \xrightarrow{\begin{pmatrix} -2 & 3 \\ -4 & 5 \end{pmatrix}} & \mathbb{R}^2 \\ \uparrow \varphi & & \uparrow \varphi^{-1} \\ \mathbb{R}^2 & & \mathbb{R}^2 \end{array}$$

$$\begin{pmatrix} -2 & 3 \\ -4 & 5 \end{pmatrix} \begin{pmatrix} 1 \\ 1 \end{pmatrix} = \begin{pmatrix} 1 \\ 1 \end{pmatrix} = 1 \cdot \begin{pmatrix} 1 \\ 1 \end{pmatrix} + 0 \cdot \begin{pmatrix} 0 \\ 1 \end{pmatrix}$$

$$\begin{pmatrix} -2 & 3 \\ -4 & 5 \end{pmatrix} \begin{pmatrix} 0 \\ 1 \end{pmatrix} = \begin{pmatrix} 3 \\ 5 \end{pmatrix} = 3 \cdot \begin{pmatrix} 1 \\ 1 \end{pmatrix} + 2 \cdot \begin{pmatrix} 0 \\ 1 \end{pmatrix}$$

$\begin{pmatrix} 1 & 3 \\ 0 & 2 \end{pmatrix}$ \leftarrow upper-triangular!

Diagonalizability: a special case (5.C)

Previously we've discussed eigenvalues and eigenvectors.

Rem: If v_1 and v_2 are eigenvectors for the same eigenvalue λ , then so is any linear combination $k_1 v_1 + k_2 v_2$.

Def: The eigenspace associated to an eigenvalue λ is $E(\lambda, f) = \ker(f - \lambda)$.

Lemma: For $\lambda \neq \lambda'$, $E(\lambda) \cap E(\lambda') = \{0\}$. \square

Cor: For $\{\lambda_i\}$, the sum $E(\lambda_1) + \dots + E(\lambda_n)$ is direct, and $\dim(E(\lambda_1) + \dots + E(\lambda_n)) = \sum_i \dim E(\lambda_i)$. \square

Rem: If f is diagonalizable, then its eigenvalues are the diagonal entries, and $V = \bigoplus_j E(\lambda_j)$.

Lemma: Take V f.d., $f: V \rightarrow V$ linear with eigenvalues $\lambda_1, \dots, \lambda_m$. TFAE:

a) f is diagonalizable. b) V has a basis of eigenvectors.

c) There exist 1-dim invariant subspaces $U_j \leq V$ with $V = \bigoplus_j U_j$.

d) $V = E(\lambda_1) \oplus \dots \oplus E(\lambda_m)$. e) $\dim V = \dim E(\lambda_1) + \dots + \dim E(\lambda_m)$.

Pf: $a \Leftrightarrow b \Rightarrow c \Rightarrow b$ are all easy. $b \Rightarrow d$ by collecting invariant subspaces of like eigenvalue. $d \Rightarrow e$ by directness. To get $e \Rightarrow b$, union bases for the individual subspaces together. \square

Cor: If f has $n = \dim V$ distinct eigenvalues, then f is diagonalizable.

Ex: $\begin{pmatrix} 1 & 1 & -1 \\ -6 & 8 & -3 \\ -4 & 4 & 1 \end{pmatrix}$ has eigenvectors $\begin{pmatrix} -1 \\ 0 \\ 2 \end{pmatrix}$, $\begin{pmatrix} 1 \\ 1 \\ 0 \end{pmatrix}$, and $\begin{pmatrix} 0 \\ 1 \\ 1 \end{pmatrix}$ and eigenvalues 3, 2, and 5.

Ex: $\begin{pmatrix} 0 & 1 & 1 \\ 0 & 0 & 1 \\ 0 & 0 & 0 \end{pmatrix}$ has eigenvectors $\begin{pmatrix} 1 \\ 3 \\ 0 \end{pmatrix}$ alone and eigenvalues 0 alone.

Ex: $\begin{pmatrix} -5 & -6 & 3 \\ 3 & 4 & -3 \\ 0 & 0 & -2 \end{pmatrix}$ has eigenvalues 1 and -2 and eigenvectors $\begin{pmatrix} -1 \\ 1 \\ 0 \end{pmatrix}$ and $\begin{pmatrix} -2 \\ 1 \\ 0 \end{pmatrix}$, $\begin{pmatrix} 1 \\ 0 \\ 1 \end{pmatrix}$.

Ex: $\begin{pmatrix} 1 & 1 \\ 0 & 1 \end{pmatrix}$ has eigenvalues 1 alone and eigenvectors $\begin{pmatrix} 1 \\ 0 \end{pmatrix}$ alone.

Lingering questions: • How can we effectively compute eigenvalues and eigenvectors?

- How can we recognize (special class of) diagonal matrices?
- How can we compute (bounds on) $\dim E(\lambda)$?

Orthogonality: (6.A-B)

angle

Def: An inner product on V is a bilinear $f^u \langle -, - \rangle: V \times V \rightarrow \mathbb{R}$ or \mathbb{C} such that: (i) $\langle v, v \rangle \geq 0$ for all $v \in V$ (requires " \geq " on $\langle v, v \rangle$...),
(ii) $\langle v, v \rangle = 0$ iff $v = 0$, (iii) $\langle u+v, w \rangle = \langle u, w \rangle + \langle v, w \rangle$,
(iv) $\langle cu, v \rangle = c\langle u, v \rangle$, and (v) $\langle u, v \rangle = \overline{\langle v, u \rangle}$.

Ex ① The dot products on \mathbb{R}^n and \mathbb{C}^n .

② For numbers $c_i \geq 0$, the modified dot product

$$u \bullet w = c_1 u_1 \overline{w_1} + \dots + c_n u_n \overline{w_n}.$$

③ $\langle f, g \rangle = \int_{-1}^1 f(x) \cdot g(x) dx$ on $V = \{ \text{integrable } f^u [-1, 1] \rightarrow \mathbb{R} \}$.

length

Def: The norm is defined by $\|v\| = \sqrt{\langle v, v \rangle}$. It satisfies $\|v\| = 0$ iff $v = 0$ and $\|\lambda v\| = |\lambda| \cdot \|v\|$.

Def: u and v are orthogonal when $\langle u, v \rangle = 0$.

Cor: If u and v are orthogonal, then $\|u+v\|^2 = \|u\|^2 + \|v\|^2$.

Pf: $\langle u+v, u+v \rangle = \langle u, u \rangle + \langle u, v \rangle + \langle v, u \rangle + \langle v, v \rangle$. \square

Thm (Cauchy-Schwarz): $|\langle u, v \rangle| \leq \|u\| \cdot \|v\|$, maximized only when $u = k \cdot v$ for a scalar k .

Pf: Write $u = \underbrace{\frac{\langle u, v \rangle}{\|v\|^2} \cdot v}_{\text{collinear with } v} + \underbrace{\left(u - \frac{\langle u, v \rangle}{\|v\|^2} \cdot v\right)}_{\text{orthogonal to } v =: w}$.

Pythag: $\|u\|^2 = \frac{|\langle u, v \rangle|^2}{\|v\|^4} \cdot \|v\|^2 + \|w\|^2 \geq \frac{|\langle u, v \rangle|^2}{\|v\|^2}$. \square

Cor (Triangle Ineq): $\|u+v\| \leq \|u\| + \|v\|$ for all u, v . \square

Pf: $\langle u+v, u+v \rangle = \langle u, u \rangle + \langle v, v \rangle + \langle u, v \rangle + \overline{\langle u, v \rangle}$

$$= \|u\|^2 + \|v\|^2 + 2\operatorname{Re} \langle u, v \rangle \leq 2|\langle u, v \rangle| \leq 2\|u\| \cdot \|v\| \\ = (\|u\| + \|v\|)^2. \quad \square$$

Def: An orthonormal basis is a basis (v_1, \dots, v_n) with $\|v_j\|=1$ and $\langle v_i, v_j \rangle = 0$ if $i \neq j$.

Lemma: In an orthonormal basis, $v = \langle v, v_1 \rangle v_1 + \dots + \langle v, v_n \rangle v_n$. \square

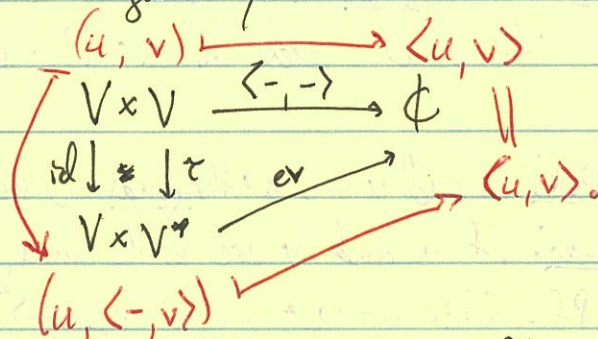
Pf: Certainly $v = k_1 v_1 + \dots + k_n v_n$. We can calculate k_j by applying $\langle -, v_j \rangle$. \square

Thm (Gram-Schmidt): Every basis can be made orthonormal.

Pf: We induct. Make v_j orthogonal to the ones before it by replacing it by $v_j - \langle v_j, v_1 \rangle v_1 - \dots - \langle v_j, v_{j-1} \rangle v_{j-1}$ and normalize that by replacing it with $v_j / \|v_j\|$. The span is preserved. \square

Rem: This procedure preserves upper-triangularity.

Recall that we have a diagram



Thm (Riesz): For V finite dim^d and $\mathcal{Q} \in V^*$, there is a unique $w \in V$ such that $\mathcal{Q} = \tau(w)$. (That is, $\langle -, - \rangle$ induces an iso $\omega: V \xrightarrow{\tau} V^*$.)

Pf: $\mathcal{Q}(v) = \mathcal{Q}(\langle v, v_1 \rangle v_1 + \dots + \langle v, v_n \rangle v_n)$
 $= \langle v, v_1 \rangle \mathcal{Q}(v_1) + \dots + \langle v, v_n \rangle \mathcal{Q}(v_n)$
 $= \langle v, \underbrace{\mathcal{Q}(v_1) v_1 + \dots + \mathcal{Q}(v_n) v_n}_{w} \rangle.$

If w_1 and w_2 both do the job, then $\langle v, w_1 \rangle - \langle v, w_2 \rangle = \mathcal{Q}(v) - \mathcal{Q}(v) = 0$
 $\langle v, w_1 - w_2 \rangle = 0$.

So, pick $v = w_1 - w_2$ and use nondegeneracy. \square
 (Rem: We already knew τ was injective by a part lemma.)

Minimization: (6.C)

Def: For an inner-product space, the annihilator gives rise to the orthogonal subspace: $U^\perp = \{v \in V \mid \langle v, u \rangle = 0 \text{ for all } u \in U\}$

Lemma: This has a number of properties, some immediate from the connection to the annihilator:

(a) $V = U \oplus U^\perp$ for f.d. U .

(b) $\dim U^\perp = \dim V - \dim U$ for f.d. V .

(c) $(U^\perp)^\perp = U$ for f.d. U . □

Def: Part (a) gives rise to the projection operator: $P_U(v) = P_U(u + u^\perp) = u$, which discards the U^\perp component of v . For an orthonormal basis e_1, \dots, e_n of U , we have $P_U(v) = \sum_j \langle e_j, v \rangle \cdot e_j$.

Lemma: (a) $\text{ker } P_U = U^\perp$.

(b) $\text{im } P_U = U$, and $P_U|_U = \text{id}$.

(c) $P_U^2 = P_U$.

(d) $\|P_U(v)\| \leq \|v\|$. □

All of these are easy to verify. The real utility of P_U is the following:

Thm: Take $v \in V$, $U \leq V$ f.d., and $u \in U$. Then $\|v - P_U v\| \leq \|v - u\|$ (with equality only at $u = P_U(v)$).

$$\begin{aligned}
 \text{Pf: } \|v - P_U(v)\|^2 &\leq \|v - P_U(v)\|^2 + \|P_U(v) - u\|^2 \\
 &= \|v - P_U(v) + P_U(v) - u\|^2 \quad (\text{Pythag.}) \\
 &= \|v - u\|^2,
 \end{aligned}$$

Equality happens iff $\|P_U(v) - u\|^2 = 0$, or $P_U(v) = u$. \square

" $P_U(v)$ is the closest point to v in U ."

Ex: As an example, we can use this to build approximations inside of function spaces.

Set $V = C[-\pi, \pi]$

$U = \text{span}\{1, x, x^2, x^3, x^4, x^5\},$

$f = \sin(x) \in V.$

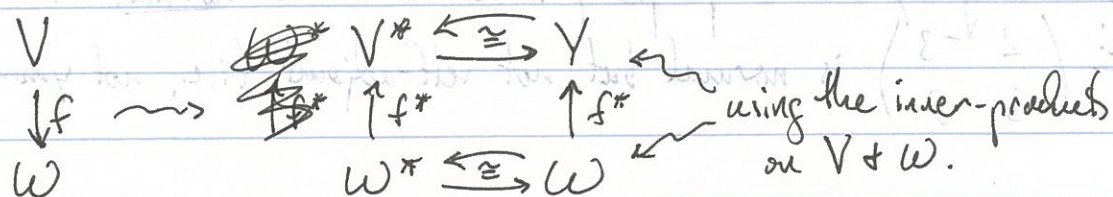
Step ①: Do Gram-Schmidt to the basis of U .

Step ②: Do orthogonal projection of f to U using this orthonormal basis.

(Timbered with Mathematica examples.)



Self-adjoint + normal operators (7.A)



So: f^* can also be considered as a map $f^*: W \rightarrow V$ in the presence of inner-products. This map is called the adjoint of f , and it satisfies $\langle f(v), w \rangle_W = \langle v, f^*(w) \rangle_V$.

Lemma: Again, we can mix annihilators + inner-products to produce:

- (a) ~~not~~ $\ker f^* = (\operatorname{im} f)^\perp$,
 - (b) $\operatorname{im} f^* = (\ker f)^\perp$,
 - (c) $\ker f = (\operatorname{im} f)^\perp$,
 - (d) $\operatorname{im} f = (\ker f)^\perp$, for $f: V \rightarrow W$.
- both f.d. \square

Lemma: For (e_1, \dots, e_n) and (f_1, \dots, f_m) orthonormal bases of V and W , the matrix M^* representing $f^*: W \rightarrow V$ is the conjugate transpose of M representing $f: V \rightarrow W$. \square

Def: ~~f is self~~ $f: V \rightarrow V$ is self-adjoint when $f = f^*$.
(Cor: In an orthonormal basis, f is conjugate-symmetric.)

These operators have particularly nice properties. Here are some:

Lemma: Every eigenvalue of a self-adjoint operator is real.

PF: $\langle Tv, v \rangle = \langle v, Tv \rangle = \langle v, \lambda v \rangle = \lambda \|v\|^2$

$\langle \lambda v, v \rangle = \lambda \|v\|^2$

\square

Lemma: Suppose V is complex and $f: V \rightarrow V$ is ~~self-adjoint~~ linear f^u .
 f is self-adjoint if + only if $\langle Tv, v \rangle \in \mathbb{R}$ for each $v \in V$. \square

Def. A slightly weaker property is for f to be normal: $f^* \circ f = f \circ f^*$.

Ex: $\begin{pmatrix} 2 & -3 \\ 3 & 2 \end{pmatrix}$ is normal but not self-adjoint (i.e., not symmetric).

Lemma: T is normal iff $\|Tv\| = \|T^*v\|$ for all v .

Pf: Consider $\langle (TT^* - T^*T)v, v \rangle = 0$. \square

Cor: T and T^* have the same eigenvectors w/ conjugate eigenvalues.

Pf: $0 = \|(T - \lambda)v\| = \|(T - \lambda)^*v\| = \|(T^* - \bar{\lambda})v\|$. \square

Lemma: If f is normal, then $v \in E(\lambda_1)$ and $w \in E(\lambda_2)$ are orthogonal.

Pf: $0 = \langle Tu, v \rangle - \langle u, T^*v \rangle = \langle \alpha u, v \rangle - \langle u, \bar{\beta} v \rangle$
 $= (\alpha - \beta) \langle u, v \rangle$. Since $\alpha - \beta \neq 0$,

we must have $\langle u, v \rangle = 0$. \square

We are going to prove the following theorem:

Thm: Let $f: V \rightarrow V$ be a linear f^* .

(a) If V is complex, then f is normal iff f is diagonalizable.

(b) If V is real, then f is self-adjoint iff f is diagonalizable.

... in an orthonormal basis.

The Spectral Theorem (7.13)

Last time, we announced two diagonalization theorems.

Today, we prove them.

Thm: V a f.d. \mathbb{C} -vector space, $f: V \rightarrow V$ linear.

f is normal iff f admits an orthonormal diagonalization.

Pf: (\Leftarrow) If f admits an orthonormal diagonalization, then f^* is diagonal for the same basis. Diagonal matrices commute.

(\Rightarrow) Start by finding an orthonormal basis in which f is upper-triangular, using Schur's theorem. We want to conclude that normality + U.T. \Rightarrow diagonal.

Write $M = \begin{pmatrix} a_{11} & a_{12} & \dots & a_{1n} \\ & a_{22} & \dots & a_{2n} \\ & & \ddots & \vdots \\ 0 & & & a_{nn} \end{pmatrix}$, so that $M^* = \begin{pmatrix} \overline{a_{11}} & \overline{a_{12}} & \dots & 0 \\ & \overline{a_{22}} & \dots & \vdots \\ & & \ddots & \vdots \\ \overline{a_{1n}} & \overline{a_{2n}} & \dots & \overline{a_{nn}} \end{pmatrix}$.

We proved last time that $\|Tv\| = \|T^*v\|$ for normal operators T , so we learn $\|a_{11}\|^2 = \|T e_1\|^2 = \|T^* e_1\|^2 = \|a_{11}\|^2 + \|a_{12}\|^2 + \dots + \|a_{1n}\|^2$.

This forces $a_{12} = \dots = a_{1n} = 0$. We can repeat this for e_2, \dots, e_n . \square

In the real case, we are much worse off: we don't even know that real operators admit U.T. presentations (or eigenvectors).

Levi: Self-adjoint real operators have eigenvectors.

Pf: Begin the same as before: starting with $v \neq 0$, find a linear dependence in $(v, f v, f^2 v, \dots, f^n v)$, guaranteed by $\dim V = n$. From the dependence $a_n f^n v + \dots + a_1 f v + a_0 v = 0$, extract a polynomial $p(x) = a_n x^n + \dots + a_1 x + a_0$, and factor it as $p(x) = c ((x-h_1)^2 + k_1^2) \dots ((x-h_m)^2 + k_m^2) (x-r_1) \dots (x-r_\ell)$. We want to show that $((f-h_j)^2 + k_j^2)$ is invertible, $k_j > 0$.

$$f^2 - 2hf + h^2 + k^2 \quad \begin{array}{l} \text{self-adj.} \\ \downarrow \end{array} \quad \begin{array}{l} \text{Cauchy-} \\ \downarrow \text{Schwarz} \end{array}$$

We just do it: $\langle (f-h)^2 + k^2 v, v \rangle = \langle f^2 v, v \rangle - 2h \langle f v, v \rangle + (h^2 + k^2) \langle v, v \rangle$
 $\geq \|f v\|^2 + 2h \|f v\| \|v\| + (h^2 + k^2) \|v\|^2$
 $= (\|f v\| + h \|v\|)^2 + (k \|v\|)^2 > 0 \text{ for } \|v\| > 0.$

With these factors eliminated, we proceed as in the \mathbb{C} case. \square

Thm. ~~\mathbb{R}~~ V a f.d. \mathbb{R} -vector space with inner product, $f: V \rightarrow V$ linear.

f is self-adjoint iff f admits an orthonormal diagonalization.

Pf. We induct on the dimension of V , since it is trivial for $\dim V = 1$.

By the Lemma, f admits an eigenvector v , spanning an invariant 1-dim^l subspace U . First, note that U^\perp is also invariant under f :

for any $u \in U$ and $v \in U^\perp$, we have $\langle u, f v \rangle = \langle f u, v \rangle = 0$, so $f v \in U^\perp$.

Additionally, $f|_{U^\perp}$ is still self-adjoint: $\langle f|_{U^\perp} v, w \rangle = \langle f v, w \rangle = \langle v, f w \rangle = \langle v, f|_{U^\perp} w \rangle$.

Hence, we can induct on $f|_{U^\perp}: U^\perp \rightarrow U^\perp$ to complete the proof. \square

That's enough for one day. To summarize:

- \mathbb{C} and normal \equiv diagonalizable orthonormally.
- \mathbb{C} and self-adjoint \equiv $\xrightarrow{\quad \text{iff} \quad}$ + real eigenvalues.
- \mathbb{R} and self-adjoint \equiv diagonalizable orthonormally.
- \mathbb{R} and normal \Leftarrow § 9.B.

Square roots and geometry (7.C)

One consequence of the spectral theorem is that, ^{because} diagonalized operators have very easy arithmetic, so do self-adjoint operators. Consider the following:

Def. An operator $f: V \rightarrow V$ is positive when it is self-adjoint and when $\langle f v, v \rangle \geq 0$ for all v . (If V is complex, we just ask for the inequality + drop the adjointness.)

Lemma. For $f: V \rightarrow V$, TFAE:

- (a) f is positive.
- (b) f is self-adjoint + all the e. values are ≥ 0 .
- (c) f has a positive square root.
- (d) f has a self-adjoint square root.
- (e) There is a second operator $g: V \rightarrow V$ with $f = g^* \circ g$.

Pf. $a \Rightarrow b$ by positivity on the e. vectors. $b \Rightarrow c$ by taking an entry-wise square root of the diagonal. $c \Rightarrow d$ trivially, as does $d \Rightarrow e$.

To get $e \Rightarrow a$, $T^* = (R^* R)^* = R^* R^{**} = R^* R = T$. \square

In fact, if we fully restrict attention to positive operators,

Lemma. ... the positive square root of a positive operator is unique.

Pf. For g a root of f and $f v = \lambda v$ an eigenvector of f , we want to show $g v = \sqrt{\lambda} v$ for g any positive square root.

We know g admits a diagonalization and that its square has eigenvalues the squares of those of g . The li. lemma for eigenvectors forces $g v = \sqrt{\lambda} v$. \square

Lingering question. How many other square roots are there?

Non-obvious defⁿ. $f: V \rightarrow V$ is an isometry if $\|f v\| = \|v\|$

for all $v \in V$. These are the "geometry-preserving f 's".

Lemma: For $f: V \rightarrow V$, TFAE:

- (a) f is an isometry.
- (b) $\langle fu, fv \rangle = \langle u, v \rangle$ for all u, v .
- (c) fe_1, \dots, fe_n is orthonormal for each orthonormal list e_1, \dots, e_n .
- (d) there exists any orthonormal list e_1, \dots, e_n s.t. fe_1, \dots, fe_n is too.
- (e) $f^* \circ f = \text{id}$.
- (f) $f \circ f^* = \text{id}$.
- (g) f^* is an isometry.
- (h) f is invertible and $f^{-1} = f^*$.

Pf: (a \Rightarrow b) This was homework: inner products can be computed from norms.

(b \Rightarrow c) Being orthonormal is an inner product condition.

(c \Rightarrow d) Trivial: pick any orthonormal basis.

(d \Rightarrow e) We have $\langle e_i, e_j \rangle = \langle f e_i, f e_j \rangle = \langle f^* f e_i, e_j \rangle$. Since $\{e_j\}$ forms a basis, this gives $\langle f^* f u, v \rangle$ for all $u, v \in V$.
This forces $f^* f = \text{id}$.

(e \Rightarrow f) Since V is finite dimensional, $f^* f = \text{id}$ forces $f f^* = \text{id}$.

(f \Rightarrow g) $\|f^* v\|^2 = \langle f^* v, f^* v \rangle = \langle f f^* v, v \rangle = \langle v, v \rangle = \|v\|^2$.

(g \Rightarrow h) Apply a \Rightarrow e n f for f^* .

(h \Rightarrow a) $\|f v\|^2 = \langle f v, f v \rangle = \langle f^* f v, v \rangle = \langle v, v \rangle = \|v\|^2$. \square

Remark: (e) is supposed to mean that there are lots of square roots of the identity, connected to the various isometries.

Polar decomposition + SVD (7.D)

Today we use our study of square roots to tackle presentations of arbitrary operators.

"Polar
decomp."

Thm: For $f: V \rightarrow V$, there is an isometry g with $f = g \circ \sqrt{f^* f}$.

Pf: Start by noting $\|fv\|^2 = \langle f^* f v, v \rangle = \langle \sqrt{f^* f} v, \sqrt{f^* f} v \rangle$ positive.

$= \|\sqrt{f^* f} v\|^2$. We "define" a function $g: \text{im } \sqrt{f^* f} \rightarrow \text{im } f$

by $g(\sqrt{f^* f}(v)) = f(v)$ — now we need to check ① that

this defⁿ is sound, ② that it extends to V , and ③ that we get

an isometry in the end.

①: $\|fv_1 - fv_2\| = \|f(v_1 - v_2)\| = \|\sqrt{f^* f}(v_1 - v_2)\| = 0$, so that
 $\ker \sqrt{f^* f} \subseteq \ker f$.

②: We also know that $\dim \text{im } f = \dim \text{im } \sqrt{f^* f}$ and that
 $\dim(\text{im } f)^\perp = \dim(\text{im } \sqrt{f^* f})^\perp$. We use this to extend g :

g acts as above on $\text{im } \sqrt{f^* f}$ and by any isometry carrying
an orthonormal basis of $(\text{im } \sqrt{f^* f})^\perp$ to $(\text{im } f)^\perp$ one of.

③ So extended, g is an isometry: g 's two definitions

are individually isometric, and the Pythagorean Theorem
extends this over the orthogonal sum. \square

Rem: Even though g and $\sqrt{f^* f}$ are diagonalizable, this may
require different orthonormal bases for each.

In fact, this isn't such a problem, because isometries are nice
enough in any basis. Favoring the orthonormal basis
diagonalizing $\sqrt{f^* f}$ leads to the Singular Value Decomposition.

Def: the singular values of f are the eigenvalues of $\sqrt{f^*f}$, with each eigenvalue repeated $\dim E(\lambda, \sqrt{f^*f})$ times. (these are the diagonal entries of an orthonormal diagonal preimage of $\sqrt{f^*f}$.)

Thm (SVD): There exist orthonormal bases (e_j) and (d_j) of V such that $f(v) = s_1 \langle v, e_1 \rangle d_1 + \dots + s_n \langle v, e_n \rangle d_n$, for (s_1, \dots, s_n) the singular values of f .

Pf: Let (e_1, \dots, e_n) present an orthonormal diagonalization of $\sqrt{f^*f}$, so that $(\sqrt{f^*f})(v) = s_1 \langle v, e_1 \rangle e_1 + \dots + s_n \langle v, e_n \rangle e_n$.

Then the columns of the isometry g appearing in the polar decomposition of f give an orthonormal set $(d_j = g(e_j))$, and $f(v) = g(s_1 \langle v, e_1 \rangle e_1 + \dots + s_n \langle v, e_n \rangle e_n)$

$$= s_1 \langle v, e_1 \rangle d_1 + \dots + s_n \langle v, e_n \rangle d_n. \quad \square$$

Rem: This is a slick, useful upgrade from Gaussian elimination, which also cleverly picked bases that diagonalized a matrix.

Rem: The e.values of $\sqrt{f^*f}$ are the nonnegative roots of the e.values of f^*f .

Ex: $f(x_1, x_2, x_3, x_4) = (0, 3x_1, 2x_2, -3x_4)$ has $f^*f(x_1, x_2, x_3, x_4) = (9x_1, 4x_2, 0, 9x_4)$, so the s.values of f are $(3, 2, 0)$, whereas the e.values of f are merely $-3 + 0$, which is not enough to recover f (since f is not normal, hence not diagonalizable).

Generalized Eigenvectors (B.A.)

In this chapter, we are aiming to correct a deficiency in our discussion of eigenspaces and diagonalization: the only operators admitting diagonal presentations are those with $V = \bigoplus E(\lambda, f)$, but in general $\bigoplus E(\lambda, f)$ may be a proper subspace of V , as with $\begin{pmatrix} 0 & 1 \\ 0 & 0 \end{pmatrix}$, which has $E(0) = \langle e_1 \rangle \subsetneq \mathbb{R}^2$. This example is instructive: the behavior of this operator stabilizes after two applications, and $E(0, f^2) = \mathbb{R}^2$ is not a proper subspace.

Lemma: For any $f: V \rightarrow V$, $0 = \ker f^0 \leq \ker f^1 \leq \ker f^2 \leq \dots$ \square

Lemma: For $N \geq \dim V$, $\ker f^N = \ker f^{N+1}$.

Pf: First, note that if $\ker f^N = \ker f^{N+1}$ is ever satisfied then $\ker f^{N+M} = \ker f^{N+M+1}$ for any $M \geq 0$: for $v \in \ker f^{N+M+1}$ we have $0 = f^{N+M+1}(v) = f^{N+1}(f^M v) \Rightarrow f^N(f^M v) = 0$. Second, we can't have an ascending chain of subspaces of length $> \dim V$. \square

Cor: For $N \geq \dim V$, $V = \ker f^N \oplus \operatorname{im} f^N$.

Pf: First check directness: $v \in (\ker f^N) \cap (\operatorname{im} f^N)$ satisfies $f^N v = 0$ and $v = f^N w$, but then $f^{2N} w = 0$ implies $f^N w = 0 = v$.

From here, apply the FTLA to $f^N: V \rightarrow V$. \square

The extreme case of this gets a special name:

Def: f is called nilpotent if $N \gg 0$ gives $\ker f^N = V$ (or $f^N = 0$).

Rem: Build a basis of $\ker f$, extend to one of $\ker f^2$, ..., etc.
 f is upper-triangular with a 0 diagonal for this basis.

This also leads us to consider the "stable" behavior of eigenvectors.

Def: The generalized eigenspace is $G(\lambda, f) = \ker (f - \lambda I)^N$, $N \geq \dim V$.

Generalized eigenvectors have properties akin to classical eigenvectors.

Lemma: If v_1, \dots, v_m are generalized eigenvectors for distinct e. values $\lambda_1, \dots, \lambda_m$, then (v_1, \dots, v_m) is a linearly independent list.

Pf: Consider a dependence $0 = a_1 v_1 + \dots + a_m v_m$. Let $k \geq 0$ be the largest value with $w = (f - \lambda_1)^k v_1 \neq 0$, so that $(f - \lambda_1)w = 0$, witnesses w as an eigenvector. Hence, we calculate

$$\begin{aligned} 0 &= (f - \lambda_1)^k (f - \lambda_2)^n \dots (f - \lambda_m)^n (a_1 v_1 + \dots + a_m v_m) \\ &= (f - \lambda_1)^k (f - \lambda_2)^n \dots (f - \lambda_m)^n (a_1 v_1) \\ &= a_1 (f - \lambda_2)^n \dots (f - \lambda_m)^n w \\ &= a_1 (\lambda_1 - \lambda_2)^n \dots (\lambda_1 - \lambda_m)^n w, \text{ which forces } a_1 = 0. \end{aligned}$$

Repeating this with other a_j gives $a_j = 0$ for each j . \square

Cor: This gives us an extension

$$\bigoplus_j E(\lambda_j, f) \leq \bigoplus_j G(\lambda_j, f) \leq V.$$

Next time: This is always an equality.

Decomposition of an operator (8.3):

Thm: For V/\mathbb{C} finite \dim , and $f: V \rightarrow V$ a linear operator, let $\lambda_1, \dots, \lambda_n$ be the eigenvalues of f . Then:

(a) $V = \bigoplus_j G(\lambda_j, f)$

(b) Each $G(\lambda_j, f)$ is invariant under f .

(c) $(f - \lambda_j) \upharpoonright_{G(\lambda_j, f)}$ is nilpotent.

Pf: (b) Note that $\ker p(f)$ and $\ker p(f)$ are invariant under f for any poly^d p .

\Rightarrow Then, $G(\lambda_j, f) = \ker (f - \lambda_j)^N$ is such a subspace.

(c) Follows from the defⁿ, since $G(\lambda_j, f) = \ker (f - \lambda_j)^N$.

(a) We induct on $\dim V$. Start by extracting an eigenvalue λ_1 of f , which decomposes V as $G(\lambda_1, f) \oplus U$, $U = \ker (f - \lambda_1)^N$, which is also an invariant subspace. We induct to get $U = G(\lambda_2, f|_U) \oplus \dots \oplus G(\lambda_m, f|_U)$ and we want to show $G(\lambda_j, f|_U) = G(\lambda_j, f)$. " \subseteq " is immediate. To get " \supseteq ", take $v \in G(\lambda_j, f)$, which we write as $v = v_1 + u$, and decompose $u = v_2 + \dots + v_m$ to get $v = v_1 + v_2 + \dots + v_m$. \uparrow $G(\lambda_1, f)$ \uparrow $G(\lambda_2, f|_U)$ \uparrow $G(\lambda_m, f|_U)$ \uparrow $G(\lambda_j, f)$
The linear independence lemma then forces $v_2 = 0$ except for v_j . In particular, $v_1 = 0$, so $v = u$, but then $v \in G(\lambda_j, f|_U)$. \square

So, if you are willing to tolerate generalized eigenvectors, you can exhaust V . Our question is then: what good is this?

Def: The algebraic multiplicity of λ_j is $\dim G(\lambda_j, f)$. The geometric multiplicity of λ is $\dim E(\lambda, f)$. (Aster just call the former "multiplicity".)

Def: Block matrices are matrices built by sewing smaller matrices together. A matrix is block diagonal if it's diagonal as a block matrix.

Cor: Every \mathbb{C} -operator admits a basis s.t. its presentation is block-diagonal with U.T. blocks.

Pf: Break V up into $\bigoplus_j G(\lambda_j, f)$. Then $f|_{G(\lambda_j, f)}$ is nilpotent, so admits an U.T. presentation w/ zeroes on the diagonal. The same basis makes $f|_{G(\lambda_i, f)}$ U.T. w/ λ_j 's on the diagonal. \square

In 8.D we will do even better than this. Right now, though, we can already find a neat application:

Lemma: If $M = I + N$ for N nilpotent, there exists \sqrt{M} .

Pf: Taylor-expand $\sqrt{1-x}$. Because N is nilpotent, we only need finitely many terms & don't care about convergence. \square

Cor: Any invertible operator $f \in \mathbb{C}$ has a square root.

Pf: Decompose f into block diagonal U.T. form. Each block can be written as $\lambda \cdot I + N$ for $\lambda \neq 0$ and N nilpotent, hence each block has a square root. Reassembling the blocks gives a square root for f . \square

Characteristic + Minimal Polynomials (B.C)

Def: The minimal polynomial of an operator f is the monic polynomial p of minimal degree such that $p(f) = 0$.

Lemma: Such a polynomial exists.

Pf: Take $n = \dim V$. Then $(1, f, f^2, \dots, f^{n^2})$ is dependent in $\mathcal{L}(V, V)$, and we take m to be the smallest index with $(1, f, \dots, f^m)$ dependent. The dependence gives a candidate polynomial. To see uniqueness, note that the difference of two candidate poly^s is another poly^d with lower degree. \square

Rem: $\deg(\text{minpoly}(f)) \leq (\dim V)^2$, by this proof.

Def: For $f \in \mathcal{L}(V)$, the characteristic polynomial of f is

$$\text{charpoly}(f) = \prod_{\lambda \text{ eigenvalue}} (z - \lambda)^{\dim G(\lambda, f)}$$

Thm (Cayley-Hamilton): The characteristic polynomial of f evaluated at f gives zero.

Pf: Decompose $V = \bigoplus_i G(\lambda_i, f)$, and shuffle the factors of $\text{charpoly}_f(z)$ so that $(z - \lambda_i)^{\dim G(\lambda_i, f)}$ appears last. This kills the vectors in $G(\lambda_i, f)$ by definition. \square

Cor: $\text{minpoly}(f) \mid \text{charpoly}(f)$, and in particular $\deg \text{minpoly} \leq \deg \text{charpoly}$.

Pf: In fact, the minimal poly divides any poly q with $q(f) = 0$.

The division algorithm gives $q = \text{min} \cdot s + r$ with $\deg r < \deg \text{min}$ and $r(f) = r(f) + \text{min}(f) \cdot s(f) = q(f) = 0$, which forces $r = 0$. \square

Our last result is that the minimality of the minimal polynomial does not remove from it the basic features of the characteristic poly.¹

Lemma: Write p for the minimal polynomial. For λ a zero of p ,
 $p(z) = (z - \lambda) \cdot q(z)$, and $p(f)(v) = (f - \lambda)q(f)(v) = 0$.
 By minimality, $q(f)(v) \neq 0$ for some v , so this is an e. vector of f with weight λ . In the other direction, if λ is an e. value of f w/ e. vector v , then $0 = p(f)(v) = p(\lambda)v$, hence $p(\lambda) = 0$. \square

Ex: $M = \begin{pmatrix} 0 & 0 & 0 & -3 \\ 1 & 0 & 0 & 6 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 0 \end{pmatrix}$, $M^2 = \begin{pmatrix} 0 & 0 & 0 & -3 & 0 \\ 0 & 0 & 0 & 6 & -3 \\ 1 & 0 & 0 & 6 \\ 0 & 1 & 0 & 0 \end{pmatrix}$, $M^3 = \begin{pmatrix} 0 & 0 & -3 & 0 & 0 \\ 0 & 0 & 6 & -3 & 0 \\ 0 & 0 & 0 & 6 & -3 \\ 1 & 0 & 0 & 0 & 6 \end{pmatrix}$

$M^4 = \begin{pmatrix} 0 & -3 & 0 & 0 & 0 \\ 0 & 6 & -3 & 0 & 0 \\ 0 & 0 & 6 & -3 & 0 \\ 0 & 0 & 0 & 6 & -3 \\ 1 & 0 & 0 & 0 & 6 \end{pmatrix}$, $M^5 = \begin{pmatrix} -3 & 0 & 0 & 0 & -18 \\ 6 & -3 & 0 & 0 & 36 \\ 0 & 6 & -3 & 0 & 0 \\ 0 & 0 & 6 & -3 & 0 \\ 0 & 0 & 0 & 6 & -3 \end{pmatrix}$ $\leadsto M^5 - 6M$ is diagonal
 $\leadsto M^5 - 6M + 3 = 0$

gives the minimal polynomial.

Fact from Math 123: There are 5 roots of this polynomial, all distinct, none expressible in terms of radical expressions. i.e., e. values!

Cor!: Computing eigenvalues exactly is not a solvable problem.

Jordan Form (8.D)

Previously we've shown the nilpotent operators admit bases in which their matrix representative is U.T. with vanishing main diagonal. Our goal today is to improve this: we will show that we can find a matrix that is nonzero only on the superdiagonal, and there has only 0's + 1's.

lem: For N nilpotent, there are vectors v_1, \dots, v_m and indices k_1, \dots, k_m such that

(a) $N^{k_1} v_1, \dots, v_1, \dots, N^{k_m} v_m, \dots, v_m$ is a basis for V .

(b) $N^{k_1+1} v_1 = \dots = N^{k_m+1} v_m = 0$.

Pf: We induct on $\dim V$, since $\dim V = 1 \Rightarrow N = 0$.

Since N is nilpotent, N is neither injective nor surjective, and we can form $N|_{\text{im } N}$. Applying the inductive hypothesis, we get ~~from~~ vectors $v_1, \dots, v_m \in \text{im } N$ and indices k_1, \dots, k_m satisfying (a) and (b). Preimage each v_j to $N(u_j) = v_j$ and trade k_j for $k_j + 1$. We claim this is at least l.i.: a dependence would image to a dependence in $\text{im } N$, leaving just $N^{k_1+1} u_1, \dots, N^{k_m+1} u_m$ unaccounted for — but there too are l.i.. Extending to a basis gives other vectors w_1, \dots, w_j with $Nw_1, \dots, Nw_j \in \text{im } N$, hence these can be perturbed to have the property $Nw_1 = \dots = Nw_j = 0$. \square

(1.8) Jordan normal form

Thm: Every \mathbb{C} -operator $f: V \rightarrow V$, $\dim V < \infty$, admit a Jordan basis, where f has a block-diagonal expression by blocks of the form $\begin{pmatrix} \lambda_i & 1 & 0 \\ & \ddots & \vdots \\ 0 & & \lambda_i \end{pmatrix}$.

Pf: Nilpotent operators were handled by the previous lemma.

In general, break up $V = \bigoplus_j G(\lambda_j, f)$, and consider $(f - \lambda_j)|_{G(\lambda_j, f)}$, which is nilpotent. A Jordan basis for $f - \lambda_j$ is also a Jordan basis for f , hence we can take the union over j . \square

Complexification (9.A)

Jordan normal form is about as much as anyone knows about nice presentations of complex operators. We now turn to real operators, where most of our theorems fail b/c we cannot assume the existence of an e-vector. Our strategy will be to replace ~~$f: V \rightarrow V$~~ $f: V \rightarrow V$ over \mathbb{R} with a complex operator that retains much of the information of f .

Def: For V/\mathbb{R} , we define $V_{\mathbb{C}}/\mathbb{C}$ by $V_{\mathbb{C}} = V \oplus iV$.

For $f: V \rightarrow V$, we define $f_{\mathbb{C}}: V_{\mathbb{C}} \rightarrow V_{\mathbb{C}}$ by $f(v + i v') = f(v) + i f(v')$.

Ex: $\mathbb{R}^n_{\mathbb{C}} \cong \mathbb{C}^n$, and thus preserves matrices:

$$\begin{array}{ccc} V \xrightarrow{f} V & \xrightarrow{\cong} & V_{\mathbb{C}} \xrightarrow{f_{\mathbb{C}}} V_{\mathbb{C}} \\ \cong \downarrow & & \downarrow \cong \\ \mathbb{R}^n \xrightarrow{M} \mathbb{R}^n & \xrightarrow{\cong} & \mathbb{C}^n \xrightarrow{M_{\mathbb{C}}} \mathbb{C}^n \end{array}$$

has the same entries as M .

lem: Real operators admit invariant subspaces of dim. 1 or 2.

Pf: $f_{\mathbb{C}}$ has an eigenvector: $f(u + i v) = (a + bi)(u + i v)$
 $= (au - bv) + i(bu + av)$. So, take $\mathcal{U} = \text{span}\{u, v\}$. \square

lem: The minimal poly^s of f and $f_{\mathbb{C}}$ agree.

Pf: For p the min. poly. of f , $p(f_{\mathbb{C}}) = p(f)_{\mathbb{C}} = 0$.

Conversely, if $q \in \mathbb{C}[x]$ satisfies $q(f_{\mathbb{C}}) = 0$, then

(Re $q(f) = 0$, so comparing degrees forces f 's min poly = $p_{\mathbb{C}}$. \square

Cor: For $\lambda \in \mathbb{R}$, λ is an e-value of f iff it's an e-value of $f_{\mathbb{C}}$. \square

lem: $(f_{\mathbb{C}} - \lambda I)(u + i v) = 0$ iff $(f - \bar{\lambda} I)(u - i v) = 0$. \square

Cor: $\lambda \in \mathbb{C}$ is an e-value of $f_{\mathbb{C}}$ iff $\bar{\lambda}$ is too, and their multiplicities agree. \square

Cor: Every real operator on an odd-dim^d space has an e-value. \square

Cor: The characteristic polynomial of $f|_F$ is actually real.

Pf: Remember the formula $\text{char poly}_f(z) = \prod_{\lambda} (z - \lambda)^{\dim G(\lambda, f)}$.

Our previous Cor says $\lambda \in \mathbb{C} \setminus \mathbb{R}$ and $\bar{\lambda}$ come in equal weights, and then factors collect to give

$$(z - \lambda)^{\dim G(\lambda, f)} (z - \bar{\lambda})^{\dim G(\bar{\lambda}, f)} = (z - 2\operatorname{Re}(\lambda)z + |\lambda|^2)^{\dim G(\lambda, f)} \quad \square$$

Def: The characteristic polynomial of a real operator is the characteristic polynomial of the complexification $f|_F$.

Consequences: $f: V \rightarrow V$ a real operator.

(a) $\deg \text{char poly}(f)(z) = \dim V$.

(b) real zeroes of $\text{char poly}(f)$ are real e. values of f .

(c) [Cayley-Hamilton] $\text{char poly}_f(f) = 0$.

(d) The minimum polynomial divides the char. polynomial.

(e) $\deg \text{min. poly} \leq \deg \text{char. poly}$. \square

Operators on real inner product spaces (9.8)

Goal: Understand normal real operators, the last case of the spectral-type results.

Start just with $\dim V = 2$:

Lemma: For $f: V \rightarrow V$, $\dim V = 2$, TFAE:

(a) f normal but not self-adjoint.

(b) All orthonormal bases present f as $\begin{pmatrix} a & -b \\ b & a \end{pmatrix}$ for $b \neq 0$.

(c) For some orthonormal basis, f is presented as $\begin{pmatrix} a & -b \\ b & a \end{pmatrix}$ for $b > 0$.

Pf: (a \Rightarrow b) Start with an orthonormal basis e_1, e_2 , presenting f as $\begin{pmatrix} a & c \\ b & d \end{pmatrix}$. Then $\|T e_1\|^2 = \|T^* e_1\|^2$ implies $a^2 + b^2 = a^2 + c^2$.

If $b = c$, we're self-adjoint, so $b = -c$ instead. Then,

$$f^* f = \begin{pmatrix} a^2 + b^2 & -ab + bd \\ -ab + bd & b^2 + d^2 \end{pmatrix} \text{ and } ff^* = \begin{pmatrix} a^2 + b^2 & ab - bd \\ ab - bd & b^2 + d^2 \end{pmatrix},$$

and $f^* f = ff^*$ forces $2ab = 2bd$, or $a = d$.

(b \Rightarrow c) Either (e_1, e_2) works or $(e_1, -e_2)$ does.

(c \Rightarrow d) Actually do the matrix mult. \square

Now, we will want an inductive decomposition of V . The following lemma assures us that this is a sane thing to do.

Lemma: V is a f.d. inner-product space, $f: V \rightarrow V$ normal, $U \subseteq V$ invariant.

(a) U^\perp is invariant under f .

(b) U is invariant under f^* .

(c) $(f|_U)^* = f^*|_{U^\perp}$.

(d) $f|_U$ and $f|_{U^\perp}$ are normal operators.

Pf: Begin by extending an orthonormal basis of U to one of V :

$(e_1, \dots, e_m, f_1, \dots, f_n)$.

Inside of this basis, f presents as $\begin{pmatrix} A & B \\ 0 & C \end{pmatrix}$, since U is invariant.

But: $\sum_{ij} |A_{ij}|^2 = \sum_{ij} \|f e_{ij}\|^2 = \sum_{ij} \|f^* e_{ij}\|^2 = \sum_{ij} |A_{ij}|^2 + \sum_{ij} |B_{ij}|^2$,

so $B \equiv$ the zero matrix. Invariance of U^\perp follows. For (b) + (c), the conjugate transpose of f 's matrix is again block-diagonal, which gives invariance of U under f^* and a calculation of $f^*|_{U^\perp}$. \square

Thm: For V a f -d. \mathbb{R} -inner product space, $f: V \rightarrow V$ is normal iff V has an orthonormal basis where f presents as block-diagonal with 1×1 scaling blocks and 2×2 scale + rotate blocks.

Pf: (\Leftarrow) Scaling + rotations all commute.

(\Rightarrow) Induct on $\dim V$. f has an invariant subspace U of $\dim 1$ or 2 , and U^\perp is invariant under f . We did the 2×2 case at the beginning. \square

*Cor: For V as above, $f: V \rightarrow V$ is an isometry iff f admits an orthonormal presentation as a block diagonal matrix as above without scaling (i.e., the 1×1 blocks are ± 1 + 2×2 's are $\begin{bmatrix} \cos \theta & -\sin \theta \\ \sin \theta & \cos \theta \end{bmatrix}$).

Pf: Isometries are normal, so the Thm applies. Because f is an isometry, it can't scale anything. \square

Thm/o proof: Real operators also admit Jordan decompositions. Each Jordan block is either (i) a complex Jordan block $\begin{pmatrix} \lambda & 0 \\ 0 & \bar{\lambda} \end{pmatrix}$ or (ii) a block diagonal matrix itself with identical 2×2 blocks $\begin{bmatrix} a_i & b_i \\ -b_i & a_i \end{bmatrix}$ (describing mult. by λ) on the diagonal + 2×2 identity blocks on the block-superdiagonal. \square

Trace (10.A):

As briefly advertised earlier in the semester, some of the coefficients of the characteristic polynomial deserve special attention: the trace and the determinant. The trace is the less interesting of the two, so we treat it first.

Def: In the expansion $(z - \lambda_1) \dots (z - \lambda_n) = z^n - (\lambda_1 + \dots + \lambda_n)z^{n-1} + \dots$, the coeff of $-z^{n-1}$ is called the trace of the operator f . (It is the sum of the e-values, repeated by algebraic mult.)

Our main goal today is to show that this value is actually computable — unlike any particular e-value alone.

Def: Given a ^{square} matrix M , the trace of the matrix is the sum $\sum_{j=1}^n M_{jj}$ of its diagonal entries.

Thm: The two definitions of the trace agree when expanding f in a basis.

lem: If A and B are matrices of the same size, then $\text{tr}(AB) = \text{tr}(BA)$.

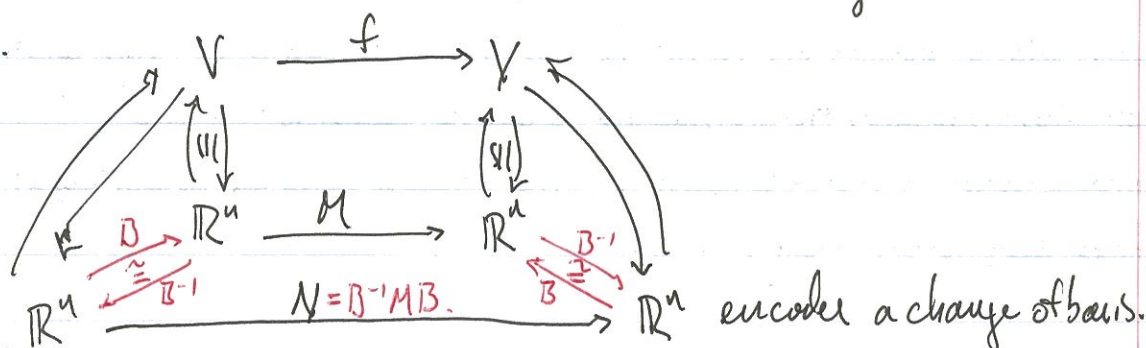
Pf: The j^{th} term on the diagonal of AB is expressed by

$$(AB)_{jj} = \sum_{k=1}^n A_{jk} B_{kj}. \text{ Summing over } j, \text{ we have}$$

$$\text{tr}(AB) = \sum_{j=1}^n (AB)_{jj} = \sum_{j=1}^n \sum_{k=1}^n A_{jk} B_{kj} = \sum_{k=1}^n \sum_{j=1}^n B_{kj} A_{jk} = \sum_{k=1}^n (BA)_{kk} = \text{tr}(BA). \quad \square$$

Cor: The trace of a matrix is invariant under change of basis.

Pf:



Hence, $\text{tr}(M) = \text{tr}(B^{-1}MB) = \text{tr}(MBB^{-1}) = \text{tr}(M)$. \square

Pf of Thm: Put f (or f_B) into upper-triangular form. There, the two definitions of trace clearly agree. Coupling this to Cor, we are done. \square

This has surprising corollaries of its own:

Cor: tr is additive: $\text{tr}(M+N) = \text{tr}(M) + \text{tr}(N)$. \square

Cor: There do not exist operators f, g with $fg - gf = \text{id}$.

Pf: $\text{tr}(fg - gf) = \text{tr}(fg) - \text{tr}(gf) = \text{tr}(fg) - \text{tr}(fg) = 0$. Meanwhile, $\text{tr}(\text{id}) = \dim V \neq 0$. \square

Determinants (10.3)

Def: The determinant of an operator f is $(-1)^{\dim V}$ times the constant coeff^t of its characteristic poly^d:

$$\text{charpoly}(z) = z^n - \text{tr}(f)z^{n-1} + \dots + (-1)^n \det(f).$$

Cor: (From homework): f is invertible iff $\det(f) \neq 0$. \square

Cor: The characteristic poly^d of f is $\det(z - f)$. ~~where f is an operator~~

Pf: Note that λ is an e-value of f iff $(z - \lambda)$ is an e-value of $z - f$:
 $-(f - \lambda) = (z - f) - (z - \lambda)$. Raising both sides to $\dim V$ and taking multiplicities also shows the algebraic multiplicities match. The characteristic poly^d of f and the determinant of $z - f$ thus match factor-wise. \square

Warning: Above, we shyly traded our k -linear map $f: V \rightarrow V$ for a $k[z]$ -linear map $f_{k[z]}: V_{k[z]} \rightarrow V_{k[z]}$ (a complexification). However, $k[z]$ is not a field! You can make this legal either by invoking modules or by using the field $k(z)$ of rat^l poly^s. You need to build gen. e-space decomposition either way, though...

As last time, we now want to start computing the determinant of operators presented as matrices. In the diagonal cases

$$\det \begin{pmatrix} a_{11} & & 0 \\ & \ddots & \\ 0 & & a_{nn} \end{pmatrix} = a_{11} \cdot a_{22} \cdot \dots \cdot a_{nn}, \text{ which is also multiplicative.}$$

However, this doesn't work for $\begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix}^2 = \begin{pmatrix} 1 & 0 \\ 0 & 1 \end{pmatrix}$.

Motivating example: $\begin{pmatrix} 0 & \dots & 0 & a_n \\ a_1 & \ddots & \ddots & 0 \end{pmatrix} = M$. Starting with v_1 , we compute $(v_1, Mv_1, M^2v_1, \dots, M^{n-1}v_1) = (v_1, a_1v_1, a_1a_2v_1, \dots, a_1 \dots a_{n-1}v_1)$, so that the l.i. of this list $\Rightarrow \deg \text{min poly} \geq n$. This forces $\text{char} = \text{min} = z^n - a_1 \dots a_n$.

So, the determinant seems to care about all diagonals, not just the main one.
We'll take this a step further:

Def: If an $n \times n$ -matrix, $\det M = \sum_{(m_1, \dots, m_n) \in \text{perm } n} \text{sign}(m_1, \dots, m_n) M_{m_1, 1} \dots M_{m_n, n}$,
where $\text{sign}(m_1, \dots, m_n)$ is -1 raised to the # of disorders in (m_1, \dots, m_n) .

Ex: $\det(a_{ij}) = a_{11}$. $\det \begin{pmatrix} a_{11} & a_{12} \\ a_{21} & a_{22} \end{pmatrix} = a_{11}a_{22} - a_{21}a_{12}$.

$\det \begin{pmatrix} a_{11} & a_{12} & a_{13} \\ a_{21} & a_{22} & a_{23} \\ a_{31} & a_{32} & a_{33} \end{pmatrix} = a_{11}a_{22}a_{33} - a_{21}a_{12}a_{33} + a_{31}a_{22}a_{13} - \dots \leftarrow 3 \text{ more terms.}$

Cor: Interchanging two columns reverses the sign of \det . Hence, if two columns are equal, $\det = 0$. \square

Lemma: For column vectors $A_1, \dots, A_{i-1}, A_{i+1}, \dots, A_n$, the map
 $(A_i \in k^n) \mapsto \det(A_1 | A_2 | \dots | A_n)$ is linear.

(That is, the determinant is a "multilinear, alternating" map.) \square

Cor: $\det(AB) = \det(A B_{1,1} | A B_{1,2} | \dots | A B_{1,n})$ (det is multiplicative.)

$$= \det \left(A \sum_{m_1=1}^n B_{m_1,1} e_{m_1} \mid A \sum_{m_2=1}^n B_{m_2,2} e_{m_2} \mid \dots \mid A \sum_{m_n=1}^n B_{m_n,n} e_{m_n} \right)$$

$$= \sum_{m_1} \dots \sum_{m_n} B_{m_1,1} \dots B_{m_n,n} \det(A e_{m_1} | \dots | A e_{m_n})$$

$$= \sum_{(m_1, \dots, m_n) \in \text{perm } n} \text{sign}(m_1, \dots, m_n) B_{m_1,1} \dots B_{m_n,n} \det A$$

$$= \sum_{\text{perm } n} \text{sign}(m_1, \dots, m_n) \det A \cdot B_{m_1,1} \dots B_{m_n,n} = \det A \cdot \det B. \quad \square$$

Cor: \det of a matrix is invariant under change of basis, and the two notions of \det agree. \square

Determinants and Volume (10.13):

Today we will investigate an important geometric aspect of determinants: their connection to volumetric properties of linear maps.

Here's the slogan for today:

Thus: For $f: \mathbb{R}^n \rightarrow \mathbb{R}^n$ a real operator on a f.d. inner product space, $\det(f)$ computes the volume of a unit cube imaged by f .

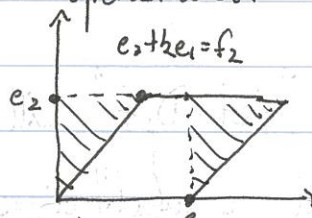
We will prove this in two ways, according to the two ways we have developed to present linear operators:

"Pf" using Gaussian elimination: We showed in a sequence of homework exercises that f can be expressed as a sequence of row and column operations applied to the identity matrix. So, we can compute $\det(f)$ by understanding the determinants of these operations.

- Scale-and-add.

These all have determinant 1.

They translate the endpoints of the parallelepiped, + this does not disturb its volume.



- Scale: These have determinant the scalar. They scale one axis of the parallelepiped, and this scales its volume by the scalar.

- Swap: These have determinant -1 using the matrix formula.

These observations collect to give a description in terms of Gr. Elimination:

identity matrix $\xrightarrow{\text{row ops}} \dots \xrightarrow{\text{row ops}} M$

unit volume $\xrightarrow{\text{row ops}} \dots \xrightarrow{\text{row ops}}$ volume of the
parallelepiped determined
by M .

Pf using Polar Decomposition. Every f factors as $f = g \circ \sqrt{f^*f}$ for some isometry g . First, note that $|\det g| = 1$, since the only eigenvalues of g satisfy $|\lambda| = 1$. Second, we know that the positive operator $\sqrt{f^*f}$ admits orthonormal eigenvalues as a diagonal matrix — whose behavior on a unit cube is easy to understand. Hence, $\det f = \det(g \circ \sqrt{f^*f}) = \det \sqrt{f^*f} = \text{vol. of unit cube under } f$. \square

Our main application of this will arise next semester.

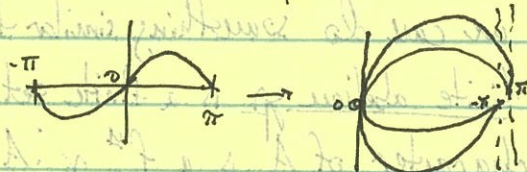
$$\begin{array}{ccc}
 M & \xrightarrow{f} & M' \\
 \downarrow \psi & & \downarrow \psi \\
 p & \xrightarrow{\quad} & f(p)
 \end{array}
 \begin{array}{l}
 \text{a smooth map} \\
 \text{derivative of } f \text{ at } p
 \end{array}
 \xrightarrow{\quad}
 \begin{array}{ccc}
 D_p f : T_p M & \longrightarrow & T_{f(p)} M' \\
 \nwarrow & & \nearrow \\
 & \text{space of tangent directions} & \\
 & \text{to } M \text{ at } p &
 \end{array}$$

The object " $\det D_p f$ " will play a role analogous to $u'(t)$ in the classical u -substitution formula $\int_{u(a)}^{u(b)} f(u) du = \int_a^b f(u(t)) \cdot u'(t) dt$.

Finite Fourier Analysis

You proved a bunch of results about $\sin nx + \cos nx$ on your hwk, considered as $f^u: [-\pi, \pi] \rightarrow \mathbb{R}$. A summary of Fourier analysis is:

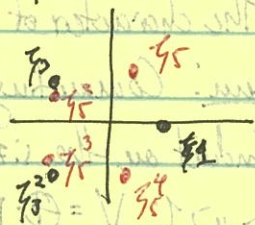
- ① There can be interpreted as f^u on the circle by gluing $-\pi$ to π .



- ② Early on, we used $e^{in\theta} = \cos(n\theta) + i\sin(n\theta)$, so we can interpret these results about \mathbb{C} -valued f^u too. In this form, $e^{i(n+m)\theta} = e^{in\theta} \cdot e^{im\theta}$.
- ③ The main Theorem: the subspace spanned by $\{e^{in\theta}\}$ is dense in all f^u , meaning any f^u can be approximated arbitrarily well using these sums.

This last theorem is beyond our reach. Today we will prove some finite analogues of it, beginning with the following:

Def: Let μ_n be the set of n^{th} roots of 1 in \mathbb{C} , i.e., $\mu_n = \{e^{2\pi i k/n} \in \mathbb{C} \mid 0 \leq k < n\}$, and let V_n be the ~~span~~ set of $f^u: \mu_n \rightarrow \mathbb{C}$.



We would like analogues of the special $f^u: e^{in\theta}$ from above, whose main property seems to be $e^{in\theta} \cdot e^{im\theta} = e^{i(n+m)\theta}$.

Def: Let $e_k \in V_n$ be the function $e_k(\zeta^k) = \zeta^{kl}$.

These satisfy $e_{n+m}(\zeta^k) = \zeta^{k(n+m)} = \zeta^{kn} \cdot \zeta^{km} = e_n(\zeta^k) \cdot e_m(\zeta^k)$.

Lemma: Under the inner product $\langle f, g \rangle = \sum_{\zeta^k \in \mu_n} f(\zeta^k) \overline{g(\zeta^k)}$, the e_k are orthogonal.

Pf: $\langle e_i, e_j \rangle = \sum_{\zeta^k \in \mu_n} \zeta^{k(i-j)} = 0$, the sum of ~~a full set of~~ roots of unity. \square

Cor: These form a basis, as they're of the right length.

Def: Given $f: \mu_n \rightarrow \mathbb{C}$, its Fourier transform \hat{f} is

$$\hat{f}(\zeta^m) = \frac{1}{n} \cdot \sum_{\zeta^k \in \mu_n} f(\zeta^k) \cdot \zeta^{-mk} = \frac{1}{n} \cdot \langle f, e_m \rangle.$$

Cor: Fourier inversion states $f(\zeta^k) = \sum_{\zeta^m \in \mu_n} \hat{f}(\zeta^m) \cdot \zeta^{mk}$.

Pf: $f = \sum_{g \in G} \hat{f}(g) \cdot e_g$, so $f(g^k) = \sum_{g \in G} \hat{f}(g) \cdot \left(\sum_{h \in G} g^{mk} \right)$. Just evaluate. \square

In fact, we can do something similar for \mathbb{C} -valued f^{un} on any finite abelian gp.

Def: A finite abelian gp is a finite set A equipped w/ a comm., unital ~~mult~~^{sum} w/ inverses.

Def: A character of A is a fⁿ $\chi: A \rightarrow \mathbb{C}$ satisfying $\chi(a+b) = \chi(a)\chi(b)$.

Lemma: Two distinct characters $\chi \neq \rho$ satisfy $\langle \chi, \rho \rangle = 0$.

Pf: Recall $\langle \chi, \rho \rangle = \frac{1}{|A|} \sum_{a \in A} \chi(a) \overline{\rho(a)} = \frac{1}{|A|} \sum_{a \in A} \chi(a) \cdot \rho^{-1}(a) = \frac{1}{|A|} \sum_{a \in A} (\chi \cdot \rho^{-1})(a)$.

We will show this sum is zero for any $\chi \neq \rho^{-1}$, w that $\chi \rho^{-1} = \eta \neq 1$.

Choose a $b \in A$ with $\eta(b) \neq 1$. Then $\eta(b) \sum_a \eta(a) = \sum_a \eta(a+b) = \sum_a \eta(a)$, so $0 = 0$. \square

Thm: The characters of A form a basis for $V_A = \{A \rightarrow \mathbb{C}\}$.

Pf: Lemma: Commuting families of unitary transformations are simultaneously diagonalizable.

Pf: Induct on the size n of the family (f_1, \dots, f_n) . $n=1$ is the spectral theorem.

For $n > 1$, $V = \bigoplus_j E(\lambda_j, f_n)$. On each eigenspace ~~f_1, \dots, f_{n-1}~~ we have $f_n f_i(v_j) = f_i f_n(v_j) = f_i(\lambda_j v_j) = \lambda_j f_i(v_j)$, so $f_i(v_j) \in E(\lambda_j, f_n)$. On $E(\lambda_j, f_n)$, any basis ~~for~~ precom $f_n|_{E(\lambda_j, f_n)}$ as $\lambda_j \cdot \text{Id}$, which commutes w/ everything. \square

Pf of Thm: Set $T_v: V_A \rightarrow V_A$ by $(T_v f)(x) = f(a+x)$. These commute, so diagonalizable them,

in a basis $(v_b) \in V_A$. Pick any such v_i — then $v_i(1) \neq 0$, since otherwise

$v(a) = (T_a v)(1) = \lambda_a v(1)$, but $\lambda_a \neq 0$. Define $w(x) = \lambda_x = v(x)/v(1)$.

We claim w is a character: $w(a+b) = (T_{a+b} w)(1) = \lambda_{a+b} w(1) = \lambda_a \lambda_b = w(a)w(b)$.

Since there are $|G|$ many v giving rise to $|G|$ many \perp w 's, we are done. \square

Cor: Set $\hat{f}(e) = \frac{1}{|A|} \sum_{a \in A} f(a) \overline{e(a)}$, for $e: A \rightarrow \mathbb{C}^*$ a character. Then

$f = \sum_{\text{characters } e} \hat{f}(e) \cdot e$, the Fourier inversion formula.

The Fast Fourier Transform + Complexity

Classical problem in complexity: sorting an unsorted list.

~~Insertion~~ "Insertion sort": Form a new list by inserting old list elements one-by-one into a new sorted list.

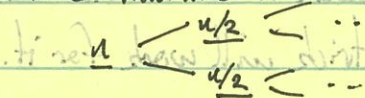
Q: How long does this take?

$(1, \dots, n) \rightsquigarrow (2, \dots, n) + (1)$
 $\rightsquigarrow (3, \dots, n) + (1, 2)$ 1 swap
 \vdots
 $1 + \dots + (n-1) = n(n-1) \cdot \frac{1}{2} = \frac{n^2}{2} - \frac{n}{2}$

Merge sort: Take a list, divide it into 2 halves, mergesort those, union them.

Q: How long does this take? Call it $f(n)$. Then $f(n) = 2f(n/2) + n$.

Picture this like:



$$n + n + \dots + n = n \cdot \lg_2 n.$$

So, mergesort is faster.

Another problem: multiplication. $\begin{array}{r} \times 5821 \\ 2642 \end{array} = 15,379,082$ requires 16 's, and some additions. In general, this algorithm takes $\approx n^2 + 2n$ steps.

Today we will use the material from last time to improve this.

Observation 1: Multiplication of integers is close to multiplication of polynomials — just with a carrying step. Namely, setting $p = 2x^3 + 6x^2 + 4x + 2$ and $q = 5x^3 + 8x^2 + 2x + 1$, we have $p(10) = 2642$, $q(10) = 5821$, and $(p \cdot q)(10) =$ the product.

Observation 2: Polynomials of degree $n-1$ are determined by their values on any n points.

In particular, picking the points μ_n from last time, $P_{n-1} \xrightarrow{a} \mu_n$ is a linear isomorphism. It's also multiplicative: $p(\zeta^k)q(\zeta^k) = (p \cdot q)(\zeta^k)$.

So, if we want to multiply two polys, we just multiply their points.

Observation 3: The map a is a form of the Fourier transform:

$$a(p)(\zeta^k) = p(\zeta^k) = \sum_{j=0}^{n-1} p_j \zeta^{jk} = (p_j, e_{-k}) = \hat{p}(k).$$

Observation 4: Computing \hat{p} is more efficient than you might think.

$$\hat{p}(k) = \sum_{j=0}^{N-1} p_j \zeta^{jk} = \sum_{j=0}^{\frac{N}{2}-1} p_{2j} (\zeta^{2k})^j + \zeta \sum_{j=0}^{\frac{N}{2}-1} p_{2j+1} (\zeta^{2k})^j$$

split into even and odd parts, which appear when computing different k . This organizes into a scheme: $p_{k+1} = p_k$, $p_{k+1/2} = p_{k+1/2}' + \zeta^{2^{k+1/2}} p_{k+1/2}'$, $p_{-1-k} = \hat{p}(k)$. This organizes into $n \lg n$ operations multiplication.

Observation 5: The Fourier inversion formula is so similar to the formula for $\hat{p}(k)$ that the same trick will work for it.

Ex: Take $p(x) = 2x^3 + 6x^2 + 4x + 2$ and $q(x) = 5x^3 + 8x^2 + 2x + 1$.

	000	001	010	011	100	101	110	111
p_{xxx}	2	4	6	2	0	0	0	0
p_{xx}	2	2	4	4	6	6	2	2
p_{x*}	8	$2+6i$	-4	$2-6i$	6	$4+2i$	2	$4-2i$
p_{lxxx}	14	$2+6i+\zeta(4+2i)$	$-4+2i$	$2-6i-\zeta(4-2i)$	2	$2+6i-\zeta(4+2i)$	$-4-2i$	$2-6i+\zeta(4+2i)$
q_{lxxx}	16	$1+8i+\zeta(2+5i)$	$-7-3i$	$1-8i-\zeta(2+5i)$	2	$1+8i-\zeta(2+5i)$	$-7+3i$	$1-8i+\zeta(2+5i)$
pg_{lxxx}	224	$-70+20i+\zeta(138+56i)$	$34-2i$	$-70-20i+\zeta(138+56i)$	4	$-70+20i+\zeta(138-56i)$	$34+2i$	$-70-20i+\zeta(138-56i)$
pg_{xxx}	2	8	30	56	72	46	10	0
	15379082							

Observation 6: In a computer implementation, we can work mod $2^N + 1$, so that 2 is an N^{th} or $2N^{\text{th}}$ root of unity ($= \zeta$), so that mult. by ζ is also fast. We can also recurse on the " pg_{xxx} " step. In all, this runs in $n \lg n \cdot \lg \lg n$ time.

Dirichlet's theorem

At the start of this class, we proved an ancient theorem:
Thm: There are infinitely many prime numbers. \square

It is easy to ask for more information than this. For instance,

Q: Are there ∞ many primes $\equiv 1 \pmod 4$? $\equiv 3 \pmod 4$?

Pf of $\equiv 3 \pmod 4$: Assume there are finitely many, and let

$(3, p_1, \dots, p_n)$ be an enumeration. Set $N = 4p_1 \cdots p_n + 3$.

Since two primes $\equiv 1 \pmod 4$ multiply to $\equiv 1 \pmod 4$, there must be a prime $\equiv 3 \pmod 4$ dividing N . Can't be 3 or p_j for any j . \square

There is no known elementary proof of the other case. There is an analytic proof, which today we will describe. The jumping-off point is:

Thm (Euler): There is a factorization $\zeta(s) = \prod_p (1 - p^{-s})^{-1}$.

Pf idea: $\frac{1}{1-p^{-s}}$ expands as $1 + \frac{1}{p^s} + \frac{1}{p^{2s}} + \dots$. Each term $\frac{1}{n^s}$ corresponds to exactly one product term by factorization into primes. \square

Thm: The series $\sum_p \frac{1}{p}$ diverges. (Note this \Rightarrow ancient theorem!)

Pf: $\log \zeta(s) = \log \sum_n \frac{1}{n^s} = \log \prod_p (1 - p^{-s})^{-1} = - \sum_p \log(1 - p^{-s})$.

The Taylor formula for \log gives $-\sum_p \log(1 - p^{-s}) = -\sum_p \left(\frac{-1}{p^s} + O\left(\frac{1}{p^{2s}}\right) \right)$
 $= \sum_p \frac{1}{p^s} + O(1)$. Finally, let $s \rightarrow 1^+$. \square

It turns out that this is the style of argument that generalizes to handle the case $p \equiv 1 \pmod 4$, — and the modification is through finite Fourier analysis. Consider the f^u $\chi: (\mathbb{Z}/4)^* \rightarrow \mathbb{C}$ defined by ~~$\chi(n) = 1$ if $n \equiv 1 \pmod 4$, -1 if $n \equiv 3 \pmod 4$~~ $\chi(1) = 1, \chi(-1) = -1$. This "extends" to all of \mathbb{Z} by $\chi(n) = \begin{cases} 0 & \text{if } n \text{ even,} \\ 1 & \text{if } n \equiv 1 \pmod 4, \\ -1 & \text{if } n \equiv 3 \pmod 4. \end{cases}$ Define $L_\chi(s) = \sum_{n=1}^{\infty} \frac{\chi(n)}{n^s} = 1 - \frac{3}{n^s} + \frac{5}{n^s} - \frac{7}{n^s} + \dots$, and $L_\chi(1) = \pi/4$. The same "pf idea" gives $L_\chi(s) = \prod_p \frac{1}{1 - \frac{\chi(p)}{p^s}}$.

Taking logs gives $\log \zeta(s) = \sum_p \chi(p) \cdot p^{-s} + O(1)$, and the fact that $\zeta(s) \rightarrow \pi/4 \neq 0$ as $s \rightarrow 1^+$ means $\sum_p \chi(p) \cdot p^{-s}$ is convergent as $s \rightarrow 1^+$. We break it into pieces:

$\sum_p \chi(p) \cdot p^{-s} = \sum_{p \equiv 1 \pmod{4}} \frac{1}{p^s} + \sum_{p \equiv 3 \pmod{4}} \frac{-1}{p^s}$. We know $\sum_p \frac{1}{p^s} \rightarrow \infty$ as $s \rightarrow 1^+$,
 \hookrightarrow adding these gives $2 \cdot \sum_{p \equiv 1 \pmod{4}} \frac{1}{p^s} \rightarrow \infty$ as $s \rightarrow 1^+$. \square

The general Theorem is:

Thm (Dirichlet): For l and q coprime, there exist ∞^{ly} many primes of the form $p = l + q \cdot k$, $k \in \mathbb{Z}$.

We're not going to prove this, but it feels a lot like the proof just given.

The main point is that $\delta_l(n) = \begin{cases} 1 & \text{if } n \equiv l \pmod{q} \\ 0 & \text{otherwise} \end{cases}$ admits a finite

Fourier expansion in terms of characters $\chi: (\mathbb{Z}/q)^* \rightarrow \mathbb{C}$, and each nontrivial such χ gives rise to a function $L_\chi(s) \rightarrow \neq 0, \neq \infty$ for $s \rightarrow 1^+$.

Once you've made it this far, you can mimic the rest of the proof above. The real meat is in the convergence of $L_\chi(1)$; we could manually calculate it, but in general this is not possible.